

# Computable thin Boolean algebras

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Let  $T$  be a theory. Its *principal extensions* are extensions of the form  $T[\varphi] = \text{Th}(T \cup \{\varphi\}) = \{\psi \text{ is a sentence} \mid T, \varphi \vdash \psi\}$  for some sentence  $\varphi$ .

Let  $T$  be a c.e. theory. We say that  $T$  *has a few c.e. extensions* if every c.e. extension  $T' \supseteq T$  has the form  $T[\varphi]$ . Note that  $T[\varphi]$  is always a c.e. extension of  $T$ . This notions was investigated in [1], where some results were proved for essentially undecidable theories with that property, and in other articles.

Below, we consider decidable theories  $T$  only. Let  $B(T)$  be the Lindenbaum-Tarski algebra of the theory  $T$  (it is a computable Boolean algebra).

Note. (1) extensions of the theory  $T$  correspond to ideals of  $B(T)$

(2) c.e. extensions of  $T$  correspond to c.e. ideals of  $B(T)$

(3)  $T$  has a few c.e. extensions  $\Leftrightarrow$  every c.e. ideal in  $B(T)$  is principal (one-generated).

We say that a computable Boolean algebra  $B$  is *thin* if every c.e. ideal in  $B$  is principal (this notion was suggested in [2]). The following theorem describes the algebraic structures of thin computable Boolean algebras.

Theorem. Let  $B$  be a computable Boolean algebra. Then the following are equivalent:

- (1)  $B \cong B'$ , where  $B'$  is a thin computable Boolean algebra;
- (2)  $B$  is an atomic Boolean algebra.

## References

- [1] Downey R.G., *Maximal theories*, Annals of Pure and Applied Logic 33, N3, 1987, 245–282.
- [2] *Handbook of Recursive Mathematics*, Elsevier, 1998.