

Symmetries of a *q*-ary Hamming Code

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| Notation | | | | |

- $\mathbb{F}_q = GF(q)$ the Galois field of order $q = p^r$
- \mathbb{F}_q^n the *n*-dimensional vector space over \mathbb{F}_q
- $d(x, y) = \#\{i: x_i \neq y_i\}$ the Hamming distance
- $w(x) = \#\{i: x_i \neq 0\}$ weight of $x \in \mathbb{F}_q^n$
- $\operatorname{supp}(x) = \{i: x_i \neq 0\}$ the support of $x \in \mathbb{F}_q^n$
- $C \subseteq \mathbb{F}_q^n$ a *q*-ary code of length *n*;
- $d(C) = \min\{d(x, y) \colon x, y \in C, x \neq y\}$ the minimum distance of *C*



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| Definitio | ons | | | |

Code equivalence

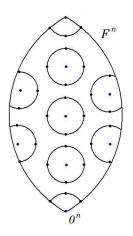
Two codes are equivalent if there is an isometry of \mathbb{F}_q^n that maps one of the codes into the other one



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Perfect codes

The balls with radius 1 centred at the codewords partition the space \mathbb{F}_q^n

Such codes have the minimum distance d = 3

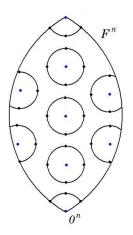
Golay codes

Binary and ternary Golay codes and codes equivalent to them have d = 7 and d = 5

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| Structure investigat | ions | | | |
| Faces of | regularity | | | |

- Linearity
- Rank of a code
- Dimension of the kernal of a code
- Perfect and uniformly packed codes
- Distance invariance
- Complete regularity
- Regular properties of minimal distance graph
- Other extremal properties



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| Structure investiga | tions | | | |
| Automo | orphisms and sy | mmetries | | |
| | | | | |

Automorphism group of a code C

The group of isometries of the space \mathbb{F}_q^n that map the code *C* into itself

Symmetry group of a code C

The group of automorphisms of *C* that fix the vector 0.



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| Isometries and auto | morphisms | | | |
| Example | es | | | |

Permutation

 $\pi \in S_n$ – a permutation on coordinate positions,

$$n = 3:$$
 $(x_1, x_2, x_3)(123) = (x_3, x_1, x_2)$

Configuration

 $\sigma = (\sigma_1, \ldots, \sigma_n) \in S_q^n - n$ permutations on elements of \mathbb{F}_q ,

$$n = 3:$$
 $(x_1, x_2, x_3)\sigma = (x_1\sigma_1, x_2\sigma_2, x_3\sigma_3)$



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| Isometries and auto | morphisms | | | |
| Isometri | es of the space | \mathbb{F}_q^n | | |

Theorem (Markov, 1956)

The group of isometries of the space \mathbb{F}_q^n is

$$\operatorname{Aut}(\mathbb{F}_q^n) = S_n \wedge S_q^n = \{(\pi; \sigma) \colon \pi \in S_n, \sigma \in S_q^n\}$$

with multiplication given by

$$(\pi;\sigma)(\tau;\delta) = (\pi\tau;\sigma\tau\cdot\delta)$$

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| Isometries and aut | omorphisms | | | |
| Kinds of | f isometries: <i>q</i> = | = 2 | | |
| | | | | |

Permutation automorphisms

 $\pi \in S_n \longrightarrow \operatorname{PAut}(\mathbb{F}_2) = \operatorname{Sym}(\mathbb{F}_2)$

All configurations are translations

| S_2 acts on \mathbb{F}_2 : | $x\in \mathbb{F}_2, \sigma\in S_2^n$ |
|--------------------------------|--------------------------------------|
| e ightarroweta+0 | $x\sigma = x + v$ |
| (01) ightarroweta+1 | for some $v \in \mathbb{F}_2$ |



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| Isometries and aut | omorphisms | | | |
| Kinds of | f isometries: <i>q</i> = | = 3 | | |
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Permutation automorphisms

 $\pi \in S_n \longrightarrow \operatorname{PAut}(\mathbb{F}_3) \subset \operatorname{Sym}(\mathbb{F}_3)$

Monomial configurations

 S_3 acts on \mathbb{F}_3 : $x \in$ $e \to 1 \cdot \beta$ $x\sigma$ $(12) \to 2\beta$ for

 $x \in \mathbb{F}_3, \sigma$ – multiplying $x\sigma = xD$ for some diagonal matrix *D*

Monomial automorphisms

 $x \in \mathbb{F}_3, \pi \in S_n, \sigma$ – multiplying $x(\pi; \sigma) = xPD = xM$ for some monomial matrix M

•
$$MAut(\mathbb{F}_3) = Sym(3)$$

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| Isometries and aut | omorphisms | | | |
| Kinds of | f isometries: $q \ge$ | <u>≥</u> 4 | | |
| | | | | |

Permutation automorphisms

 $\pi \in S_n \longrightarrow \operatorname{PAut}(\mathbb{F}_q^n) \subset \operatorname{Sym}(\mathbb{F}_q^n)$

Configurations

- q = 4Gal(\mathbb{F}_4), ×, + (0 $\alpha^2 1 \alpha$) $\rightarrow (\beta + \alpha)^2$, where α – primitive element of \mathbb{F}_4 no matrix representation for all!
- *q* ≥ 5 no field operations for all!



Introduction 000 Linear codes

The Hamming code

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Conclusion

Linear and semilinear transformations of \mathbb{F}_{q}^{n}

General linear group

 $\begin{aligned} & \operatorname{GL}_n(q) \\ & f(\alpha x + \beta y) = \alpha f(x) + \beta f(y) \\ & \text{for all } x, y \in \mathbb{F}_q^n \text{ and } \alpha, \beta \in \mathbb{F}_q \end{aligned}$

General semilinear group

$$\begin{split} &\Gamma \mathcal{L}_n(q) = \operatorname{Gal}(\mathbb{F}_q) \land \operatorname{GL}_n(q) \\ &f(\alpha x + \beta y) = \gamma(\alpha) f(x) + \gamma(\beta) f(y) \\ &\text{for all } x, y \in \mathbb{F}_q^n \text{, all } \alpha, \beta \in \mathbb{F}_q \text{, and some } \gamma \in \operatorname{Gal}(\mathbb{F}_q) \end{split}$$

But not all of them are isometries of \mathbb{F}_q^n !

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MacWilliams' theorem

Theorem (MacWilliams, 1962)

Two linear codes are monomially equivalent iff there exists an isomorphism between them (as linear spaces) preserving the weight of each vector

Corollary 1

 $\operatorname{MAut}(\mathbb{F}_q^n)$ – all linear symmetries of \mathbb{F}_q^n

Corollary 2

 $\operatorname{Gal}(\mathbb{F}_q) \prec \operatorname{MAut}(\mathbb{F}_q^n)$ – all semilinear symmetries of \mathbb{F}_q^n

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The automorphism group of a linear code

Proposition

If a code $C \subseteq \mathbb{F}_q^n$ *is linear, then*

 $\operatorname{Aut}(C) \cong \operatorname{Sym}(C) \land C$

Theorem

• *C* is an $[n, n - m, d \ge 3]_q$ -code

⇒ The semilinear symmetry group of C is isomorphic to some subgroup of $\Gamma L_m(q)$



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| What is | s known | | | | |
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Semilinear automorphisms of \mathcal{H}

Theorem

The semilinear symmetry group of *a q*-ary Hamming code \mathcal{H} *of length* $n = \frac{q^m - 1}{q - 1}$ *is isomorphic to* $\Gamma L_m(q)$

- q = 2, 3 all symmetries of \mathbb{F}_q^n are linear Aut $(\mathcal{H}) \cong \operatorname{GL}_m(q) \times \mathcal{H}$
- $q \ge 4$ not all symmetries of \mathbb{F}_q^n are semilinear Aut $(\mathcal{H}) \cong \Gamma L_m(q) \land \mathcal{H} ?$



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| What is known | | | | |

Is there anything to doubt?

Example

- $C \subset \mathbb{F}_q^n$ is the linear code with $H = [1 \ 1 \dots 1]$
- $A \in S_q$ is a linear transformation of \mathbb{F}_q as a vector space over the subfield \mathbb{F}_p
- \Rightarrow $(e, (A, A, \dots, A)) \in \operatorname{Aut}(C)$
- ⇒ for $q \ge 8$ there exists A such that this automorphism of C is neither linear nor semilinear



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| Results | | | | |
| Collinea | ar triples | | | |

•
$$T = \{x \in \mathcal{H} \colon w(x) = 3\}$$

• Sym $(\mathcal{H}) \leq$ Sym(T)

Lemma

- $x, y \in T$
- supp(x) = supp(y)
- \Rightarrow $y = \mu x$ for some $\mu \in \mathbb{F}_q^*$



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| Results | | | | |
| Symme | tries of Hammir | ng triples | | |
| | | | | |
| Lemm | a (saving collinearity | y) | | |

- $(\pi; \sigma) \in \operatorname{Sym}(T)$
- \Rightarrow (π ; σ) preserves collinearity of vectors from \mathbb{F}_q^n

Lemma (saving sum)

- $(\pi; \sigma) \in \operatorname{Sym}(T)$
- \Rightarrow $(\pi; \sigma)$ preserves sum of vectors from \mathbb{F}_q^n

Lemma

- $(\pi; \sigma) \in \operatorname{Sym}(T)$
- $\Rightarrow (\pi; \sigma)$ is a semilinear transformation of \mathbb{F}_q^n

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| Results | | | | |
| The aut | omorphism gro | up of ${\cal H}$ | | |

Theorem

For any *q*-ary Hamming code \mathcal{H} of length $n = \frac{q^{m-1}}{q-1}$, where $q, m \ge 2$, it is true

 $\operatorname{Aut}(\mathcal{H}) \cong \Gamma \operatorname{L}_m(q) \checkmark \mathcal{H}$



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| Conclusion | | |
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We proved that

- all symmetries of the Hamming code are semilinear
- the same can be said about the triple system of a *q*-ary Hamming code



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