On completion of latin hypercuboids of order 4

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Let Σ be a nonempty set (the alphabet). A function $f: \Sigma^n \to \Sigma$ is called an *n*-ary quasigroup if $f(\overline{x}) \neq f(\overline{y})$ for all $\overline{x}, \overline{y} \in \Sigma^n$ such that $d(\overline{x}, \overline{y}) = 1$ where $d(\overline{x}, \overline{y})$ is the Hamming distance.

A function defined on $\Omega \subset \Sigma^n$ and satisfying the same condition is called a partial *n*-ary quasigroup. We deal only with partial *n*-ary quasigroups $f : \Sigma^{n-1} \times \Sigma' \to \Sigma$, where $\Sigma' \subset \Sigma$.

The table of values of an *n*-ary quasigroup is called a *latin hypercube* (the *n*-dimensional generalization of latin square) and the table of values of a partial *n*-ary quasigroup is called a *latin hypercuboid*.

A subset $M \subset \Sigma^{n+1}$ is called an *MDS code (with distance 2)* if $|M| = |\Sigma|^n$ and $d(\overline{x}, \overline{y}) \ge 2$ for all distinct $\overline{x}, \overline{y} \in M$. It is clear that the graph of an *n*-ary quasigroup *f* is an MDS code. Moreover there exists a one-to-one correspondence between the *n*-ary quasigroups and the MDS codes (with distance 2).



There are some constructions of perfect codes (with distance 3) that use MDS codes (with distance 2).

Zinov'ev V. A. Generalized cascade codes. (Russian) Problemy Peredachi Informacii **12**, no. 1, 5–15. 1976.

Phelps K. T. A general product construction for error correcting codes. *SIAM J. Algebraic Discrete Methods.* **5**, no. 2, 224–228. 1984.

Moreover, MDS codes are used for classifying perfect codes with rank +2 with respect to the linear perfect code.

Avgustinovich S. V., Heden O., Solov'eva F. I. The classification of some perfect codes. *Des. Codes Cryptogr.* **31**, no. 3, 313–318. 2004.

A partial *n*-ary quasigroup *f* is *extendable* if $f = q|_{\Sigma^{n-1} \times \Sigma'}$, where *q* is an *n*-ary quasigroup.

By the definitions a latin hypercuboid can be completed to a latin hypercube if and only if the partial *n*-ary quasigroup is extendable.





The fact that every latin rectangle can be completed to a latin square is a simple consequence of Konig's theorem. It is well known that every partial *n*-ary quasigroup $f : \Sigma^{n-1} \times \Sigma' \to \Sigma$ is extendable if $|\Sigma'| = 1$ or $|\Sigma'| = |\Sigma| - 1$. Therefore, all partial *n*-ary quasigroups of order $M = |\Sigma| \leq 3$ are extendable.

Kochol proved that for all k and M satisfying M/2 < k < M - 2 there is an $M \times M \times k$ latin cuboid that cannot be completed to a latin cube.

Kochol M. Relatively narrow latin parallelepipeds that cannot be extended to a latin cube. Ars Comb., **40**, 247–260, 1995. Kochol M. Latin $(n \times n \times (n-2))$ -parallelepipeds not completing to a latin cube. Math. Slovaka, **39**, 3–9, 1989.

Examples of non completable $M \times M \times k$ latin cuboids for M = 5, 6, 7, 8 and k = 2, 2, 3, 4 respectively are constructed by McKay and Wanless.

McKay B. D., Wanless I. M. A census of small Latin hypercubes. *SIAM J. Discrete Math.*, **22**, (2), 719–736, 2008.

Using these facts it is easy to show that there are *n*-dimensional latin hypercuboids of order *M* that cannot be completed to a latin hypercube for $n \ge 3$ and $M \ge 5$. However the question of existence of a non completable latin hypercuboid for M = 4 was open.

Main result

Theorem. Every latin hypercuboid of order 4 is completable to a latin hypercube.

As mentioned above we need to complete only $4 \times \cdots \times 4 \times 2$ latin hypercuboids because the other cases, $4 \times \cdots \times 4 \times 1$ and $4 \times \cdots \times 4 \times 3$, are trivial. In other words, we have a pair

of (n-1)-ary quasigroups f_0 and f_1 such that $f_0(\overline{x}) \neq f_1(\overline{x})$ for all $\overline{x} \in \Sigma^{n-1}$ and we want to extend the pair to an *n*-ary quasigroup.

The proof of the theorem is based on the classification theorem of n-ary quasigroups of order 4.

Theorem (classification of *n*-ary quasigroups of order 4). Every *n*-ary quasigroup of order 4 is permutably reducible or semilinear.

Krotov D. S. and Potapov V. N. *n*-Ary quasigroups of order 4. *SIAM J. Discrete Math.*, **23** (2), 561–570, 2009.

An *n*-quasigroup *f* is termed *permutably reducible* if there exist $m \in \{2, ..., n-1\}$, an (n - m + 1)-quasigroup *h*, an *m*-quasigroup *g*, and a permutation $\sigma : \{1, ..., n\} \rightarrow \{1, ..., n\}$ such that

$$f(x_1,\ldots,x_n) \equiv h(g(x_{\sigma(1)},\ldots,x_{\sigma(m)}),x_{\sigma(m+1)},\ldots,x_{\sigma(n)}).$$

Let $\Sigma = \{0, 1, 2, 3\}$. Consider the 2-fold MDS code (the union of two disjoint MDS codes) $L \subset \Sigma^n$ defined by the indicator function

$$\chi_L(x_1,\ldots,x_n)\equiv\chi_{0,1}(x_1)\oplus\cdots\oplus\chi_{0,1}(x_n).$$

We say that a 2-MDS code is *linear* if it is isotopic to the 2-MDS code *L*. Let *f* be an *n*-ary quasigroup and let $a, b \in \Sigma, a \neq b$. Then the set

$$S_{a,b}(f) = \{(x_1,\ldots,x_n) \in \Sigma^n \mid f(x_1,\ldots,x_n) \in \{a,b\}\}$$

is a 2-fold MDS code. An *n*-ary quasigroup f is called *semilinear* if for some $a, b \in \Sigma$ the set $S_{a,b}(f)$ is linear.



By analyzing the various cases of types of pairs of (n-1)-ary quasigroups f_0 and f_1 , we can prove the theorem.