Triply even codes obtained from some graphs and finite geometries

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$L(K_6) = T(6) = J(6, 2)$ $\binom{6}{2} = 15$ SRG($v = 15, k = 8, \lambda = 4, \mu = 4$)





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- In general, for T(n),
 - Tonchev (1988), Brouwer-van Eijl (1992): dim C,
 - Haemers, Peeters and van Rijckevorsel (1999): weight distribution

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We say that a vector $\boldsymbol{x} \in \mathbb{F}_2^n$ is

even
$$\iff$$
 wt(x) $\equiv 0 \pmod{2}$
doubly even \iff wt(x) $\equiv 0 \pmod{4}$
triply even \iff wt(x) $\equiv 0 \pmod{8}$

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Even, doubly even, and triply even codes

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(i)
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 is doubly even for all $i \in \{1, ..., k\}$,
(ii) wt $(r_i * r_j) \equiv 0 \pmod{2}$ for all $i, j \in \{1, ..., k\}$.
(denoting by * the entrywise product)

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This is because the mapping

$$f: \{ even vectors in \mathbb{F}_2^n \} \to \mathbb{F}_2$$

defined by

$$f: \mathbf{x} \mapsto rac{\mathsf{wt}(\mathbf{x})}{2} \mod 2$$

is a quadratic form.

$$f(\sum_{i=1}^k a_i r_i) = \sum_{i=1}^k a_i^2 f(r_i) + \sum_{i < j} a_i a_j \operatorname{wt}(r_i * r_j).$$

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Do these property imply C is triply even? No, in general. We need: (iii) wt $(r_h * r_i * r_j) \equiv 0 \pmod{2}$ for all $h, i, j \in \{1, ..., k\}$. The number of common neighbors of three vertices $\equiv 0 \pmod{2}$

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Code of T(4n + 2) is triply even

The code C generated by the row vectors of the adjacency matrix of T(4n+2) is triply even.

The code *C* generated by the row vectors of the adjacency matrix of T(4n+2) is triply even. dim C = n - 2. Note dim C = n - 2

$$rate = rac{\dim C}{\operatorname{length}} = rac{n-2}{\binom{4n+2}{2}} o 0 \quad (n \to \infty).$$

Image: Image:

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Theorem (Betsumiya–M., 2012)

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Sharp contrast with doubly even codes: rate is always $\approx 1/2$. There are triply even codes with rate 1/4 whenenver $n \equiv 0 \pmod{16}$.

Classification of triply even codes of length 48

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Motivation comes from Framed Vertex Operator Algebras (FVOA).

- The moonshine module V^t has Virasoro frames, and each Virasoro frame gives rise to a triply even code of length 48.
- Lam-Yamauchi (2008) showed that, conversely, every triply even code of length divisible by 16 is obtained from some FVOA.
- The classification lead Lam and Shimakura to discover new FVOA≈CFT conjectured by Schellekens (1993).

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$$\{\pm x\} \sim \{\pm y\}$$

$$\iff Q(x \pm y) = 0$$

$$\iff \text{the line through } \langle x \rangle \text{ and } \langle y \rangle \text{ is a "tangent"}$$

$$\iff \text{the line } \langle x, y \rangle \text{ and the surface } Q = 0 \text{ in } PG(3, 3)$$

$$\text{has exactly one point in common}$$

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Let $V = \mathbb{F}_q^4$ be equipped with a nondegenerate quadratic form with Witt index 1, for example, with $\eta \notin (\mathbb{F}_q^{\times})^2$,

$$Q(x_1, x_2, x_3, x_4) = x_1 x_2 + x_3^2 - \eta x_4^2,$$

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Then $|X| = q(q^2 + 1)/2$. Adjacency by tangent. Not SRG unless q = 3. Brouwer–Cohen–Neumaier, Section 12.2 shows this is a 3-class association scheme.

Theorem (Betsumiya–M.)

For any odd prime power q, the code of this graph is triply even, of dimension at least $(q^2 - 1)/2$.

Proof: k, λ, μ (BCN, Section 12.2)

Let $V = \mathbb{F}_q^4$ be equipped with a nondegenerate quadratic form Q with Witt index 1. Define a graph Γ whose vertex set is

$$X = \{\{\pm x\} \mid Q(x) = 1\},\$$

with adjacency

$$\{\pm x\} \sim \{\pm y\} \iff B(x,y) = \pm 1,$$

where

$$B(x,y) = \frac{1}{2}(Q(x+y) - Q(x) - Q(y)).$$

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The graph has valency $q^2 - 1 \equiv 0 \pmod{8}$, $\lambda = 2(q - 1) \equiv 0 \pmod{4}$, $\mu = 2(q - 1)$ or $2(q + 1) \equiv 0 \pmod{4}$, depending on $\langle x, y \rangle$ is external or secant.

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Let $\langle x_1 \rangle, \langle x_2 \rangle, \langle x_3 \rangle$ be distinct vertices, Their common neighbors are $\{\langle z \rangle \mid Q(z) = 1, B(x_i, z) = \pm 1 \ (i = 1, 2, 3)\}.$

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For simplicity, assume $W = \langle x_1, x_2, x_3 \rangle$ is a nondegenerate 3-dimensional subspace, and consider

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$$\{\langle x_0 + y \rangle \mid Q(x_0 + y) = 1, y \in W^{\perp}\}$$

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We claim Γ has induced $\frac{q+1}{2}K_{q-1}$. Write $V = V_+ \oplus V_-$, where dim $V_{\pm} = 2$, V_+ contains a nonzero vector x with Q(x) = 0, V_- is anisotropic. The following subset of vertices induces $\frac{q+1}{2}K_{q-1}$:

$$egin{aligned} &Y=\{\langle\lambda x+y
angle\mid\lambda\in\mathbb{F}_q^{ imes},\;y\in V_-,\;Q(y)=1\}\ &=\{\langle\lambda x+y_i
angle\mid\lambda\in\mathbb{F}_q^{ imes},\;1\leq i\leq (q+1)/2\}, \end{aligned}$$

since

$$B(\lambda x + y_i, \mu x + y_j) = B(y_i, y_j) = \begin{cases} 1 & \text{if } i = j, \\ \text{not } \pm 1 & \text{otherwise.} \end{cases}$$

Maximality?

Since dim $C \ge (q^2 - 1)/2$, the rate is at least

$$rac{rac{q^2-1}{2}}{rac{q(q^2+1)}{2}} = rac{q^2-1}{q(q^2+1)} o 0 \quad (q o \infty).$$

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Thank you for your attention!