

Association schemes, Schur rings and the isomorphism problem for circulant graphs. Part 1.

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- Coherent configurations and coherent (cellular) algebras.
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- Schur rings and Cayley graphs isomorphism problem.
- Theory of Schur rings.

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Lecture 1.

Notation: binary relations

Let $R, S \subseteq \Omega^2$ be binary relations. Then

- $S^* := \{(\alpha, \beta) : (\beta, \alpha) \in S\}$;
- S is **symmetric** (**antisymmetric**) if $S = S^*$ ($S \cap S^* = \emptyset$ resp.);
- $\alpha S := \{\beta \in \Omega : (\alpha, \beta) \in S\}$, $S\alpha := \alpha S^*$;
- $D(S) := \{\alpha \in \Omega : \alpha S \neq \emptyset\}$, $R(S) := D(S^*)$;
- $RS = \{(\alpha, \beta) : \alpha R \cap S\beta \neq \emptyset\}$;
- $R^+ = \bigcup_{i=0}^{\infty} R^i$ is the transitive closure of R ;
- $1_{\Omega} := \{(\omega, \omega) : \omega \in \Omega\}$.

Each permutation $g \in \text{Sym}(\Omega)$ is considered as a binary relation. Thus $\alpha g = \{\alpha^g\}$ and $g^* = g^{-1}$.

Notation: partitions

- $\mathcal{P} \vdash \Omega$ means that \mathcal{P} is a partition of Ω .
- $\mathcal{P} \sqsubseteq \mathcal{C} \iff \mathcal{C}$ is a refinement of \mathcal{P} ;
- Lattice operations are denoted as $\mathcal{P} \vee \mathcal{C}$ and $\mathcal{P} \wedge \mathcal{C}$ ($\mathcal{P} \vee \mathcal{C}$ is the partition the classes of which are the intersections of classes of \mathcal{P} and \mathcal{C});
- if $\mathcal{P} \vdash \Omega$ then \mathcal{P}^{\cup} denotes the set of all possible unions of elements in \mathcal{P} ;
- $\mathcal{C} \vdash \Omega^2 \implies \mathcal{C}^* := \{C^* : C \in \mathcal{C}\}$.

Algorithms

- an algorithm can be thought as a sequence of **steps** from what we have (**input**) to what we want to get (**output**),
- depending on a computational model one has to define the **size** of the input/output, and possible "elementary" steps,
- the **complexity** of an algorithm is the number of "elementary" steps as the function of the input size.

Example: transitivity test

Input: a group $G \leq \text{Sym}(\Omega)$ given by a generator set S .

Output: "YES" or "NO" depending on whether G is transitive.

Step 1. Take $\alpha \in \Omega$ and set $\Omega_0 := \{\alpha\}$. Set $S := S \cup \{1_G\}$.

Step 2. While $(\Omega_0)^S \neq \Omega_0$ do $\Omega_0 := (\Omega_0)^S$.

Step 3. Output "YES" if $\Omega_0 = \Omega$, and "NO" otherwise.

Exercise: The complexity of the algorithm is $O(nm)$ where $n = |\Omega|$ and $m = |S|$.

Graphs

In what follows **graph** is a pair $\Gamma = (\Omega, E)$ where Ω is a finite set of **vertices** and $E \subset \Omega \times \Omega$ is the set of (directed) **edges/arcs**.

Definition.

Graphs $\Gamma_1 = (\Omega_1, E_1)$ and $\Gamma_2 = (\Omega_2, E_2)$ are called **isomorphic**, $\Gamma_1 \cong \Gamma_2$, if there is a bijection $f : \Omega_1 \rightarrow \Omega_2$ such that

$$\forall \alpha_1, \beta_1 \in \Omega_1 : (\alpha_1^f, \beta_1^f) \in E_2 \iff (\alpha_1, \beta_1) \in E_1.$$

Such a bijection is called the **isomorphism** from Γ_1 to Γ_2 ; the set of all of them is denoted by $\text{Iso}(\Gamma_1, \Gamma_2)$. The **automorphism group** of Γ is defined to be $\text{Aut}(\Gamma) := \text{Iso}(\Gamma, \Gamma)$.

The Graph Isomorphism Problem (ISO)

Problem ISO

Input: graphs Γ_1 and Γ_2

Output: “YES” or “NO” depending on whether or not $\Gamma_1 \cong \Gamma_2$.

- Given n -vertex graphs Γ_1 and Γ_2 , and a bijection $f : \Omega_1 \rightarrow \Omega_2$ one can test in time $O(n^2)$ whether $f \in \text{Iso}(\Gamma_1, \Gamma_2)$.
- Therefore $\text{ISO} \in \text{NP}$.
- An exhaustive search of all the possible bijections runs in exponential time $O(n!)$.
- At present it is not known whether $\text{ISO} \in \text{P}$.

The proof of the time bound of the best algorithm (up to now) for the ISO depends on the [Classification of Finite Simple Groups](#).

Theorem (L.Babai, E.Luks and W.Kantor, 1984)

The isomorphism of n -vertex graphs can be tested in time $\exp(O(\sqrt{n \log n}))$.

Some problems equivalent to the ISO

The problems equivalent to the ISO (Mathon,1979):

- **IMAP:** given Γ and Γ' find $f \in \text{Iso}(\Gamma, \Gamma')$ (if it exists),
- **ICOUNT:** given Γ and Γ' find $|\text{Iso}(\Gamma, \Gamma')|$,
- **ACOUNT:** given Γ find $|\text{Aut}(\Gamma)|$,
- **AGEN:** given Γ find a generators of the group $\text{Aut}(\Gamma)$,
- **APART:** given Γ find its [automorphic partition](#) $\text{Orb}(\text{Aut}(\Gamma))$.

Isomorphism problem for colored graphs

Definition

A triple (Ω, Y, c) where $c : \Omega^2 \rightarrow Y$ is a surjection, is called a [colored graph](#) with the [coloring function](#) c .

The [color classes](#) form a partition $\mathcal{C} := \{c^{-1}(y) \mid y \in Y\}$ of Ω^2 .

Definition

Two colored graphs (Ω, Y, c) and (Δ, Z, d) are [isomorphic](#) if there exists a bijection $f : \Omega \rightarrow \Delta$ such that

$$d(\alpha^f, \beta^f) = c(\alpha, \beta).$$

Remarks

- The ISO for colored graphs is equivalent to the ISO.
- One can consider a colored graph also as a partition.

Configurations and their isomorphisms

Definition.

A pair (Ω, \mathcal{C}) where $\mathcal{C} \vdash \Omega^2$ is called a [configuration](#) on Ω . The elements of \mathcal{C} are called [basic relations](#) or [color classes](#).

Let $\mathcal{X} = (\Omega, \mathcal{C})$ and $\mathcal{X}' = (\Omega', \mathcal{C}')$ be two configurations. Then

- $\text{Iso}(\mathcal{X}, \mathcal{X}') := \{f : \Omega \rightarrow \Omega' : f \text{ is a bijection and } \mathcal{C}^f = \mathcal{C}'\}$;
- $\mathcal{X} \cong \mathcal{X}' \iff \text{Iso}(\mathcal{X}, \mathcal{X}') \neq \emptyset$;
- each $f \in \text{Iso}(\mathcal{X}, \mathcal{X}')$ induces a bijection f^* between \mathcal{C} and \mathcal{C}' via $\mathcal{C}^{f^*} = \mathcal{C}'^f$, $\mathcal{C} \in \mathcal{C}$;
- $\text{Iso}(\mathcal{X}) := \text{Iso}(\mathcal{X}, \mathcal{X}) = \{g \in \text{Sym}(\Omega) : \mathcal{C}^g = \mathcal{C}\}$;
- $\text{Aut}(\mathcal{X}) := \{g \in \text{Sym}(\Omega) : \forall C \in \mathcal{C} : C^g = C\}$.

Cayley graphs and their isomorphisms

Definition

A **Cayley graph** over a finite group H defined by a **connection set** $S \subseteq H$ has H as a set of nodes and arc set

$$\text{Cay}(H, S) := \{(x, y) \mid xy^{-1} \in S\}.$$

A **circulant** graph is a Cayley graph over a cyclic group.

Definition

Two Cayley graphs $\text{Cay}(H, S)$ and $\text{Cay}(K, T)$ are **Cayley isomorphic** if there exists a group isomorphism $f : H \rightarrow K$ which is a graph isomorphism too, that is

$$\text{Cay}(H, S)^f = \text{Cay}(K, T) \iff S^f = T.$$

The graphs $\text{Cay}(\{\pm 1\}, \mathbb{Z}_5)$ and $\text{Cay}(\{\pm 2\}, \mathbb{Z}_5)$ are Cayley isomorphic.

A characterization of Cayley graphs

The automorphism group of a Cayley graph $\text{Cay}(H, S)$ contains a regular subgroup $H_R \leq \text{Sym}(H)$ consisting the right translations

$$h_R : x \mapsto xh$$

for all $h \in H$.

Theorem (Sabidussi)

A graph $\Gamma = (\Omega, E)$ is isomorphic to a Cayley graph over a group H if and only if $\text{Aut}(\Gamma)$ contains a regular subgroup isomorphic to H .

Proof.

- Pick a base point $\omega \in \Omega$.
- Define a bijection $f : \Omega \rightarrow H$, $\alpha \mapsto \bar{\alpha}$ where $\omega^{\bar{\alpha}} = \alpha$.
- Set $S = \{\bar{\alpha} : \alpha \in E\omega\}$.
- Now $f \in \text{Iso}(\Gamma, \text{Cay}(H, S))$:

$$(\alpha, \beta) \in E \iff (\omega^{\bar{\alpha}}, \omega^{\bar{\beta}}) \in E \iff (\omega^{\bar{\alpha}\bar{\beta}^{-1}}, \omega) \in E \iff \bar{\alpha}\bar{\beta}^{-1} \in S.$$

Cayley representations of graphs

Definition

Under a **Cayley representation** of a graph Γ over a group H we mean any graph $\text{Cay}(H, S)$ isomorphic to Γ .

Remark: graph can have Cayley representations over non-isomorphic groups.

Exercise:

- What happens with S if one replaces the base point?
- Find all Cayley representations of $\text{Cay}(\mathbb{Z}_6, \{\pm 1, 3\})$ and 3-dim cube.

Lecture 2.

Isomorphism problems for Cayley graphs

Problems

Given a group H and graphs $\Gamma = \text{Cay}(H, S)$ and $\Gamma' = \text{Cay}(H, S')$ we have the following problems:

- **IMAP:** find $f \in \text{Iso}(\Gamma, \Gamma')$ (if it exists),
- **ICOUNT:** find $|\text{Iso}(\Gamma, \Gamma')|$,
- **ACOUNT:** find $|\text{Aut}(\Gamma)|$,
- **AGEN:** find generators of the group $\text{Aut}(\Gamma)$.

Recognizing problem:

CGR: given a graph Θ test whether it is a Cayley graph over H .

Isomorphism problem for finite groups

Construction

For a finite group K set $\Gamma(K)$ to be a graph with vertex set $K \times K$ such that

$$(a, b) \sim (c, d) \iff a = c \vee b = d \vee ab = cd.$$

Theorem (Moorhouse, 1991)

$$K_1 \cong K_2 \iff \Gamma(K_1) \cong \Gamma(K_2).$$

Exercise

- $\Gamma(K)$ is a Cayley graph over $K \times K$.
- $\Gamma(\mathbb{Z}_4) \not\cong \Gamma(\mathbb{Z}_2 \times \mathbb{Z}_2)$:

$$\begin{array}{c} \mathbb{Z}_4 \rightarrow \\ \begin{array}{cccc} & 0 & 1 & 2 & 3 \\ 1 & 2 & 3 & 0 & \\ 2 & 3 & 0 & 1 & \\ 3 & 0 & 1 & 2 & \end{array} \end{array} \qquad \begin{array}{c} \mathbb{Z}_2 \times \mathbb{Z}_2 \rightarrow \\ \begin{array}{cccc} & 0 & 1 & 2 & 3 \\ 1 & 0 & 3 & 2 & \\ 2 & 3 & 0 & 1 & \\ 3 & 2 & 1 & 0 & \end{array} \end{array}$$

Naive vertex classification

Denote by $d_\Gamma(\alpha)$ the **valency** of the vertex α in the graph Γ , and by $d_\Gamma(\alpha, C)$ the valency of α in a vertex color class C .

Naive vertex classification

Given a graph Γ we construct an $\text{Aut}(\Gamma)$ -invariant vertex coloring.

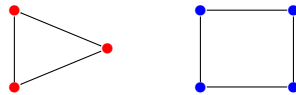
- Put vertices α and β in the same class iff $d_\Gamma(\alpha) = d_\Gamma(\beta)$;
- Iteratively, put vertices α and β in the same class iff $d_\Gamma(\alpha, C) = d_\Gamma(\beta, C)$ for all color classes C .

Comments

- The algorithm correctly finds $\text{Orb}(\text{Aut}(\Gamma))$ for the class of trees (G.Tinhofer, 1985), for almost all graphs (L.Babai, P.Erdős, S.Selkow, 1980).
- The algorithm fails when Γ is regular and $\text{Aut}(\Gamma)$ is intransitive.

The Weisfeiler-Leman algorithm

No automorphism moves red points to blue ones:



To distinguish vertices we need to color edges of Γ .

WL-algorithm (Weisfeiler-Leman (1968))

Input: graph $\Gamma = (\Omega, E)$.

Output: an $\text{Aut}(\Gamma)$ -invariant coloring of Ω^2 .

Step 1. Set $\mathcal{C} = \{1_\Omega\} \cup \{E\} \cup \{\Omega^2 \setminus (E \cup 1_\Omega)\}$.

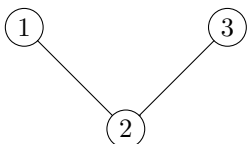
Step 2. For all $(\alpha, \beta) \in \Omega \times \Omega$ and $R, S \in \mathcal{C}$ find the number

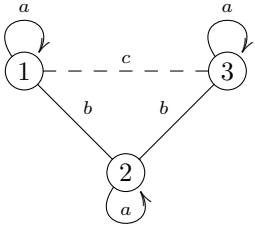
$$c(\alpha, \beta; R, S) = |\alpha R \cap S \beta|.$$

Step 3. Build a new partition $\mathcal{C} := \text{bl}(\mathcal{C})$ by putting (α, β) and (α', β') to the same class if $|\alpha R \cap S \beta| = |\alpha' R \cap S \beta'|$ for all $R, S \in \mathcal{C}$.

Step 4. Repeat the procedure till $|\mathcal{C}|$ stop to increase.

The WL-algorithm: very small example





Initial coloring

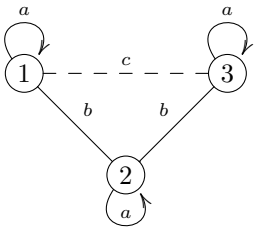
Graph:

Adjacency matrix:

$$A = \begin{pmatrix} a & b & c \\ b & a & b \\ c & b & a \end{pmatrix}.$$

First iteration

Graph:

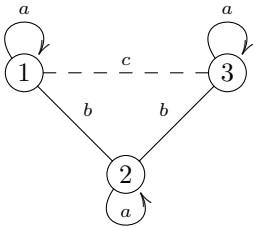


Adjacency matrix:

$$A^2 = \begin{pmatrix} a^2 + b^2 + c^2 & ab + ba + cb & ac + b^2 + ca \\ ba + ab + bc & 2b^2 + a^2 & bc + ab + ba \\ ca + b^2 + ac & cb + ba + ab & c^2 + b^2 + a^2 \end{pmatrix}.$$

First iteration

Graph:



Adjacency matrix:

$$A^2 = \begin{pmatrix} a^2 + b^2 + c^2 & ab + ba + cb & ac + b^2 + ca \\ bc + ab + ba & 2b^2 + a^2 & bc + ab + ba \\ ac + b^2 + ca & ab + ba + cb & c^2 + b^2 + a^2 \end{pmatrix}.$$

Second iteration

Graph:

Adjacency matrix:

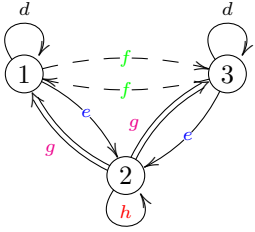
$$A = \begin{pmatrix} d & e & f \\ g & h & g \\ f & e & d \end{pmatrix}$$

WL-refinement (operation \mathfrak{bl})

The correctness of the WL-algorithm (i.e. that it is stopped) follows from the following statements.

Lemma

1. $\mathcal{C} \sqsubseteq \mathcal{S} \implies \mathfrak{bl}(\mathcal{C}) \sqsubseteq \mathfrak{bl}(\mathcal{S})$,
2. $\mathcal{C}^* = \mathcal{C} \implies \mathfrak{bl}(\mathcal{C})^* = \mathfrak{bl}(\mathcal{C})$,
3. $1_\Omega \in \mathcal{C}^\cup \implies \mathcal{C} \sqsubseteq \mathfrak{bl}(\mathcal{C})$.



Theorem

If a partition $\mathcal{C} \vdash \Omega^2$ satisfies the conditions $1_\Omega \in \mathcal{C}$ and $\mathcal{C}^* = \mathcal{C}$, then $\text{bl}(\mathcal{C})$ satisfies the same conditions and $\mathcal{C} \sqsubseteq \text{bl}(\mathcal{C})$.

Canonical WL-refinement

Proposition

Let $f : \Omega \rightarrow \Delta$ be a bijection that maps a partition \mathcal{C} of Ω^2 onto a partition \mathcal{T} of Δ^2 (i.e. $\mathcal{C}^f = \mathcal{T}$). Then $\text{bl}(\mathcal{C})^f = \text{bl}(\mathcal{T})$.

Corollary

Let Γ be a graph and \mathcal{C} the stable partition obtained by the WL-algorithm. Then $\text{Aut}(\Gamma) = \text{Aut}(\mathcal{C})$.

Given an ordered partition $\vec{\mathcal{C}} = (S_1, \dots, S_m)$ of Ω^2 the WL-algorithm produces a unique (**canonical**) ordering of the refinement $\text{bl}(\mathcal{C})$ (denoted as $\text{bl}(\vec{\mathcal{C}})$), i.e.

$$\vec{\mathcal{C}}^f = \vec{\mathcal{T}} \implies \text{bl}(\vec{\mathcal{C}})^f = \text{bl}(\vec{\mathcal{T}})$$

Coherent configurations (D. Higman, 1970)

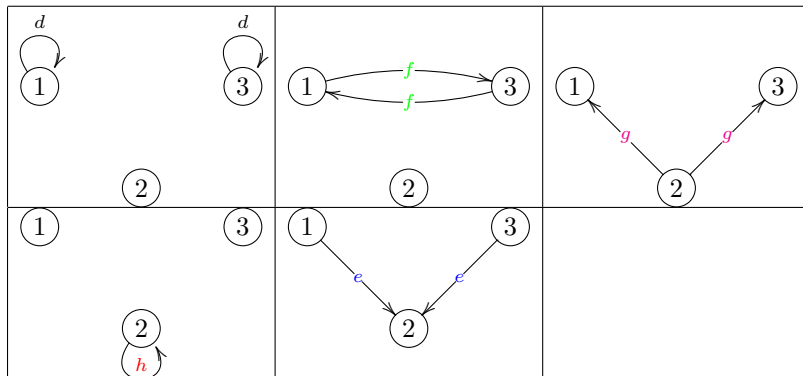
The output partition of the Weisfeiler-Leman algorithm forms a **coherent configuration**, i.e. a pair $\mathcal{X} = (\Omega, \mathcal{C})$ such that:

1. \mathcal{C} is a partition of $\Omega \times \Omega$,
2. $1_\Omega \in \mathcal{C}^\cup$,
3. $\mathcal{C}^* = \mathcal{C}$,
4. $\mathcal{C} = \text{bl}(\mathcal{C})$, or, equivalently, for all $R, S, T \in \mathcal{C}$ the **intersection number** $c_{RS}^T = |\alpha R \cap S \beta|$ does not depend on the choice of $(\alpha, \beta) \in T$.

Terminology:

- the **degree** and **rank** of \mathcal{X} are the numbers $|\Omega|$ and $|\mathcal{C}|$,
- the **basic relations** and **relations** of \mathcal{X} are the elements of \mathcal{C} and of \mathcal{C}^\cup ,
- the configuration \mathcal{X} is **homogeneous** (or **association scheme**, or **scheme**), if $1_\Omega \in \mathcal{C}$.

Coherent configurations: a concrete example.



Coherent configurations: relations

Proposition

Let $\mathcal{X} = (\Omega, \mathcal{C})$ be a coherent configuration. Then

- the set \mathcal{C}^\cup is closed with respect to boolean operations;
- $1_\Omega, \Omega^2 \in \mathcal{C}^\cup$;
- $(\mathcal{C}^\cup)^* = \mathcal{C}^\cup$;
- \mathcal{C}^\cup is closed with respect to relational product.

Coherent configurations: fibers

Definition

A **fiber** of \mathcal{X} is a set $\Delta \subset \Omega$ such that $1_\Delta \in \mathcal{C}$; the set of all fibers is denoted by $\Phi = \Phi(\mathcal{X})$.

Proposition

- \mathcal{X} is a scheme if and only if $|\Phi| = 1$,
- $\Omega = \bigcup_{\Delta \in \Phi} \Delta$,
- for any $S \in \mathcal{C}$ the sets $D(S)$ and $R(S)$ are fibers of \mathcal{X} ,
- for any $S \in \mathcal{C}$ and $\alpha \in D(S)$ we have $|\alpha S| = c_{S S^*}^T$ where $T = 1_{D(S)}$; the number $n_S = c_{S S^*}^T$ is called the **valency** of S ,
- for any $\Delta \in \Phi$ the set $\mathcal{C}_\Delta := \{C \in \mathcal{C} : D(C) = \Delta, R(C) = \Delta\}$ form a homogeneous cc on Δ , called a **homogeneous constituent** of \mathcal{C} .

Exercises

Exercises:

- Prove that if the **symmetrization** $\{C \cup C^* \mid C \in \mathcal{C}\}$ of a cc $\mathcal{X} = (\Omega, \mathcal{C})$ is a cc then \mathcal{C} is a scheme.
- Enumerate up to an isomorphism all cc of degree 2, 3, 4.
- Prove that for any triple $R, S, T \in \mathcal{C}$ it holds that

$$c_{RS}^{T^*} |T^*| = c_{ST}^{R^*} |R^*| = c_{TR}^{S^*} |S^*|.$$

Isomorphisms between coherent configurations

Definition

Two coherent configuration $\mathcal{X} = (\Omega, \mathcal{C})$ and $\mathcal{X}' = (\Omega', \mathcal{C}')$ are called **(combinatorially) isomorphic** if there exist bijections $f : \Omega \rightarrow \Omega'$ and $\phi : \mathcal{C} \rightarrow \mathcal{C}'$ such that

$$\forall_{\alpha, \beta \in \Omega} (\alpha, \beta) \in C \iff (\alpha^f, \beta^f) \in C^\phi.$$

The set of all isomorphisms between \mathcal{X} and \mathcal{X}' is denoted as $\text{Iso}(\mathcal{X}, \mathcal{X}')$.

Comments

- ϕ is induced by f , and hence is usually omitted;
- $\text{Iso}(\mathcal{X}) := \text{Iso}(\mathcal{X}, \mathcal{X})$;
- $\text{Aut}(\mathcal{X}) = \{f \in \text{Sym}(\Omega) : S^f = S \text{ for all } S \in \mathcal{C}\}$,
- $\text{Aut}(\mathcal{X}) \trianglelefteq \text{Iso}(\mathcal{X})$.

Lecture 3.

Coherent configurations generated by a graph

Theorem

Let $\langle\langle \Gamma \rangle\rangle \vdash \Omega^2$ be the **WL-closure** of a graph $\Gamma = (\Omega, E)$ obtained by applying WL-algorithm to Γ . Then

1. $E \in \langle\langle \Gamma \rangle\rangle^\cup$;
2. $\text{Aut}(\Gamma) = \text{Aut}(\langle\langle \Gamma \rangle\rangle)$.

Examples: strongly regular graphs

Definition

An undirected graph $\Gamma = (\Omega, E)$ is called **strongly regular** if its WL-closure has rank three. In other words, WL-algorithm stops at the first iteration and $\langle\langle\Gamma\rangle\rangle = \{1_\Omega, E, E^c\}$.

Proposition

A graph $\Gamma = (\Omega, E)$ is strongly regular if and only if there exist non-negative integers k, λ, μ such that

1. Γ is k -regular,
2. any two adjacent vertices have exactly λ common neighbors,
3. any two non-adjacent vertices have exactly μ common neighbors.

Exercises

Let $\Gamma = (\Omega, E)$ be a strongly regular graph with parameters v, k, λ, μ . Then

- the complement to Γ is strongly regular;
- $(k - \lambda - 1)k = (v - k - 1)\mu$;
- Γ is disconnected iff $\lambda = k - 1$ iff Γ is a disjoint union of K_k ;
- the graph $\Gamma(K)$ built from the Cayley table of a group K is strongly regular;
- the complements to $\Gamma(\mathbb{Z}_4)$ and $\Gamma(\mathbb{Z}_2 \times \mathbb{Z}_2)$ are the only strongly regular graph with parameters $v = 16, k = 6$ and $\lambda = \mu = 2$.

Examples: permutation groups

Let $G \leq \text{Sym}(\Omega)$ be a permutation group. It acts on $\Omega \times \Omega$:

$$(\alpha, \beta)^g := (\alpha^g, \beta^g), \quad \alpha, \beta \in \Omega, \quad g \in G.$$

Proposition

Set $\mathcal{C} := \text{Orb}(G, \Omega \times \Omega)$. Then

1. $\text{Inv}(G) := (\Omega, \mathcal{C})$ is a coherent configuration (of G),
2. the basic relations of \mathcal{X} are the 2-orbits of G ,
3. $\Phi(\mathcal{X}) = \text{Orb}(G, \Omega)$,
4. \mathcal{X} is a scheme if and only if G is transitive.

Definition.

A coherent configuration \mathcal{X} is **schurian** if $\mathcal{X} = \text{Inv}(G)$ for some G .

The **Schurity problem** is to check whether or not given a coherent configuration is schurian.

Exercise

Let $G = \mathbb{Z}_2^3$ and $G_1 = \langle(1, 0, 0)\rangle, G_2 = \langle(0, 1, 0)\rangle, G_3 = \langle(0, 0, 1)\rangle$.

Set $\Omega = G/G_1 \cup G/G_2 \cup G/G_3$ and $\mathcal{C} = \text{Orb}(G, \Omega^2)$.

- what is the rank of $\mathcal{X} = (\Omega, \mathcal{C})$?
- find the fibers and describe the basic graphs of \mathcal{X} ;
- find homogeneous constituents;
- find $\text{Iso}(\mathcal{X})$ and $\text{Aut}(\mathcal{X})$;

Galois correspondence between cc and groups

Definition

Let $\mathcal{X} = (\Omega, \mathcal{C}), \mathcal{X}' = (\Omega, \mathcal{C}')$ be two coherent configurations. We say that \mathcal{X} is a **fusion** of \mathcal{X}' (equivalently \mathcal{X}' is a **fission** of \mathcal{X}), notation $\mathcal{X} \sqsubseteq \mathcal{X}'$ if $\mathcal{C} \sqsubseteq \mathcal{C}'$.

Proposition

Let $\mathcal{X}, \mathcal{X}'$ be two coherent configurations on Ω and $G, H \leq \text{Sym}(\Omega)$. Then

- $\mathcal{X} \sqsubseteq \mathcal{X}' \implies \text{Aut}(\mathcal{X}) \geq \text{Aut}(\mathcal{X}')$;
- $H \leq G \implies \text{Inv}(H) \supseteq \text{Inv}(G)$;
- $G \leq \text{Aut}(\text{Inv}(G))$ and $\mathcal{X} \sqsubseteq \text{Inv}(\text{Aut}(\mathcal{X}))$.

Thus the mappings $\mathcal{X} \mapsto \text{Aut}(\mathcal{X})$ and $G \mapsto \text{Inv}(G)$ form a Galois correspondence between coherent configurations on Ω and permutation groups on Ω .

Galois closed objects

Definition

- The group $G^{(2)} := \text{Aut}(\text{Inv}(G))$ is called a **2-closure** of $G \leq \text{Sym}(\Omega)$; G is **2-closed** if $G = G^{(2)}$.
- The cc $\text{Sch}(\mathcal{X}) := \text{Inv}(\text{Aut}(\mathcal{X}))$ is called a **schurian closure** of a cc $\mathcal{X} = (\Omega, \mathcal{C})$; \mathcal{X} is **schurian** if $\text{Sch}(\mathcal{X}) = \mathcal{X}$.

Theorem

The mappings (Aut, Inv) are bijections between schurian coherent configurations on Ω and 2-closed subgroups of $\text{Sym}(\Omega)$.

The ISO is polynomially equivalent to the problem of finding the schurian closure of a coherent configuration.

Matrix operations

Notation

Let $A, B \in M_\Omega(\mathbb{F})$ be arbitrary matrices. We denote by

- AB (or $A \cdot B$) the usual matrix product;
- $A \circ B$ the Schur-Hadamard (component-wise) product, i.e. $(A \circ B)_{\alpha\beta} := A_{\alpha\beta} B_{\alpha\beta}$;
- A^\top the transposed of A ;
- I_Ω the identity matrix;
- J_Ω the all one matrix.

Proposition

The algebra $(M_\Omega(\mathbb{F}), \circ)$ is a commutative associative algebra with identity J_Ω . It is isomorphic to \mathbb{F}^n where $n = |\Omega|^2$.

Idempotents

Let (\mathcal{A}, \star) be finite dimensional algebra over field \mathbb{F} .

Definition

- $e \in \mathcal{A}, e \neq 0$ is **\star -idempotent** if $e \star e = e$;
- idempotents e, f are **orthogonal** if $e \star f = f \star e = 0$;
- idempotent e is **minimal** if it is not a sum $e = e_1 + e_2$ of pairwise orthogonal idempotents e_1, e_2 .

Exercise:

- a matrix $E \in M_\Omega(\mathbb{F})$ is \cdot -idempotent iff E is similar to a $(0, 1)$ -diagonal matrix;
- a matrix $E \in M_\Omega(\mathbb{F})$ is \circ -idempotent iff E is $(0, 1)$ -matrix iff $E = A(S)$ is the **adjacency** matrix of some $S \subseteq \Omega^2$;
- a symmetric matrix is \cdot - and \circ -idempotent iff it's $(0, 1)$ -diagonal matrix.

Coherent (cellular) algebras

Definition.

A subspace $\mathcal{A} \leq M_\Omega(\mathbb{F})$ is called a **coherent** (or **cellular**) algebra if it contains I_Ω, J_Ω and is closed with respect to \cdot, \circ, \top .

Examples: $\langle I_\Omega, J_\Omega \rangle$ and $M_\Omega(\mathbb{F})$.

Proposition

Let $\mathcal{X} = (\Omega, \mathcal{C})$ be a coherent configuration. Then the linear span $\mathbb{F}\mathcal{X} = \mathbb{F}\mathcal{C} := \langle A(C) \rangle_{C \in \mathcal{C}}$ is a coherent algebra of dimension $|\mathcal{C}|$. It is called the **adjacency** (or **Bose-Mesner**) algebra of \mathcal{X} .

The basis $\{A(C) : C \in \mathcal{C}\}$ is called the **standard** basis of $\mathbb{F}\mathcal{X}$; it consists of $(0, 1)$ -matrices and they are minimal \circ -idempotents. Moreover,

$$A(S)A(T) = \sum_{R \in \mathcal{C}} c_{ST}^R A(R), \quad S, T \in \mathcal{C}.$$

Coherent algebras are adjacency algebras

Theorem

Every coherent algebra $\mathcal{A} \leq M_\Omega(\mathbb{F})$ has a unique basis consisting of minimal \circ -idempotents which are pairwise orthogonal. If $\text{char}(\mathbb{F}) = 0$ then \mathcal{A} is the adjacency algebra of a uniquely determined coherent configuration.

Proof. \mathcal{A} is a \circ -subalgebra of $(M_\Omega(\mathbb{F}), \circ) \cong (\mathbb{F}^n, \circ), n = |\Omega|^2$.

Lemma

Let \mathcal{A} be a k -dimensional subalgebra of (\mathbb{F}^n, \circ) . Then there exists a unique basis A_1, \dots, A_k of \mathcal{A} consisting of minimal, pairwise orthogonal, \circ -idempotents. Each \circ -idempotent of \mathcal{A} is a $(0, 1)$ -linear combination of A_1, \dots, A_k and $A_1 + \dots + A_k$ is the unit of \mathcal{A} .

Proof of the Theorem (the end)

- $A_i = A(R_i), R_i \subseteq \Omega^2$;
- $i \neq j \implies A_i \circ A_j = 0 \implies R_i \cap R_j = \emptyset$;
- $\sum_{i=1}^k A_i = J_\Omega \implies \bigcup_i R_i = \Omega^2$;
- $I_\Omega = \sum_i A_i \implies 1_\Omega = \bigcup_i R_i$;
- $A_i A_j = \sum_k c_{ij}^k A_k$ for some $c_{ij}^k \in \mathbb{F}$;
- $(A_i A_j)_{\alpha\beta} = c_{ij}^k$ where k is defined by $(\alpha, \beta) \in R_k$;
- if $\text{char}(\mathbb{F}) = 0$, then $(A_i A_j)_{\alpha\beta} = |\alpha R_i \cap R_j \beta| \implies |\alpha R_i \cap R_j \beta| = c_{ij}^k$.

Thus coherent configurations and coherent algebras are the same objects presented by different languages. In what follows we'll switch between the languages freely.

Lecture 4.

Isomorphisms between coherent algebras

Definition

Given two coherent algebras $\mathcal{A} \leq M_\Omega(\mathbb{F}), \mathcal{A}' \leq M_{\Omega'}(\mathbb{F})$, a linear bijection $L : \mathcal{A} \rightarrow \mathcal{A}'$ is called an **(algebraic) isomorphism** if

- $L(XY) = L(X)L(Y)$;
- $L(X \circ Y) = L(X) \circ L(Y)$;
- $L(X^\top) = L(X)^\top$.

Proposition

Let $L : \mathbb{F}\mathcal{X} \rightarrow \mathbb{F}\mathcal{X}'$ be an algebraic isomorphism between the adjacency algebras of coherent configurations \mathcal{X} and \mathcal{X}' . Then there exists a bijection $\varphi : \mathcal{C} \rightarrow \mathcal{C}'$ such that $L(A(C)) = A(C^\varphi)$ and $c_{RS}^T = c_{R^\varphi S^\varphi}^{T^\varphi}$. Vice versa, any such φ extends uniquely up to an algebraic isomorphism between $\mathbb{F}\mathcal{X}$ and $\mathbb{F}\mathcal{X}'$; φ is called an **algebraic isomorphism** from \mathcal{X} onto \mathcal{X}' .

Properties of algebraic isomorphisms

Proposition

Let φ be an algebraic isomorphism from a coherent configuration $\mathcal{X} = (\Omega, \mathcal{C})$ to a coherent configuration $\mathcal{X}' = (\Omega', \mathcal{C}')$. Then

- $(RS)^\varphi = R^\varphi S^\varphi$ for any $R, S \in \mathcal{C}^\cup$;
- for each $\Delta \in \Phi(\mathcal{X})$ there exists a unique $\Delta' \in \Phi(\mathcal{X}')$ such that $(1_\Delta)^\varphi = 1_{\Delta'}$;
- the mapping $\Delta \mapsto \Delta'$ is a bijection from $\Phi(\mathcal{X})$ onto $\Phi(\mathcal{X}')$;
- $D(S^\varphi) = D(S)^\varphi$ and $R(S^\varphi) = R(S)^\varphi$ for all $S \in \mathcal{C}$;
- $|n_{S^\varphi}| = |n_S|$ for each $S \in \mathcal{C}$;
- $|\Delta^\varphi| = |\Delta|$ for any $\Delta \in \Phi(\mathcal{X})$.

Isomorphisms between coherent algebras

Proposition

Given $f \in \text{Iso}(\mathcal{X}, \mathcal{X}')$ the mapping $\varphi_f : S \mapsto S^f$ is an algebraic isomorphism from \mathcal{X} onto \mathcal{X}' ; we say that φ_f is induced by the isomorphism f .

All algebraic automorphisms of \mathcal{X} to itself form a subgroup of $\text{Sym}(\mathcal{C})$ denoted as $\text{Alg}(\mathcal{X})$. Notice that

$$\text{Iso}(\mathcal{X}) / \text{Aut}(\mathcal{X}) \hookrightarrow \text{Alg}(\mathcal{X}).$$

Proposition

Given $A \leq \text{Alg}(\mathcal{X})$ the subspace $(\mathbb{F}\mathcal{X})^A := \{x \in \mathbb{F}\mathcal{X} : x^a = x \text{ for all } a \in A\}$

is a coherent algebra; the corresponding coherent configuration is denoted by \mathcal{X}^A and is called an [algebraic fusion](#) of \mathcal{X} .

Coherent closure: definition

Proposition

Let $\mathcal{X} = (\Omega, \mathcal{C})$ and $\mathcal{X}' = (\Omega, \mathcal{C}')$ be coherent configurations. Then

- $\mathbb{F}\mathcal{X} \subseteq \mathbb{F}\mathcal{X}'$ if and only if $\mathcal{C} \subseteq \mathcal{C}'$;
- $\mathbb{F}\mathcal{X} \cap \mathbb{F}\mathcal{X}' = \mathbb{F}(\mathcal{C} \wedge \mathcal{C}')$ is a coherent algebra.

In general, $\mathbb{F}(\mathcal{C} \vee \mathcal{C}')$ is not necessarily coherent algebra.

Proposition

Let $A_1, \dots, A_m \in M_\Omega(\mathbb{F})$. Then the intersection of all coherent algebras containing A_1, \dots, A_m is a coherent algebra; it is called the [coherent closure](#) of A_1, \dots, A_m and denoted by $\langle\langle A_1, \dots, A_m \rangle\rangle$.

The definition of the coherent closure is not constructive; the constructive one is given by means of the WL-algorithm.

Coherent closure: properties

Definition

Given matrices $A_1, \dots, A_k \in M_\Omega(\mathbb{F})$ we define a partition $\mathcal{P}(A_1, \dots, A_k)$ of Ω^2 via the following equivalence relation:

$$(\alpha, \beta) \sim (\gamma, \delta) \iff (A_i)_{\alpha\beta} = (A_i)_{\gamma\delta} \text{ for all } 1 \leq i \leq k.$$

Proposition

Let $\mathcal{X} = (\Omega, \mathcal{C})$ be a coherent configuration. Then

1. if $A_1, \dots, A_k \in \mathbb{F}\mathcal{X}$, then $\mathcal{P}(A_1, \dots, A_k) \subseteq \mathcal{C}$,
2. if $\mathcal{S} \vdash \Omega^2$ and $\mathcal{S} \subseteq \mathcal{C}$, then $\text{bl}(\mathcal{S}) \subseteq \mathcal{C}$.

Coherent closure: computation

Finding coherent closure by means of the WL-algorithm:

Given matrices $A_1, \dots, A_k \in M_\Omega(\mathbb{F})$ do:

Step 1. Set $\mathcal{S}_0 := \mathcal{P}(A_1, \dots, A_k, A_1^\top, \dots, A_k^\top, I_\Omega)$ and $i := 0$;

Step 2. While $\mathcal{S}_i \neq \text{bl}(\mathcal{S}_i)$ do $i := i + 1$ and $\mathcal{S}_i := \text{bl}(\mathcal{S}_i)$;

Step 3. Output $\mathbb{F}\mathcal{S}$ where $\mathcal{S} = \mathcal{S}_i$.

Theorem

The above algorithm computes the coherent closure $\langle\langle A_1, \dots, A_k \rangle\rangle$ of the matrices A_1, \dots, A_k .

Proof. Let $\mathcal{X} = (\Omega, \mathcal{C})$ be the underlying coherent configuration of $\langle\langle A_1, \dots, A_k \rangle\rangle$, i.e. $\langle\langle A_1, \dots, A_k \rangle\rangle = \mathbb{F}\mathcal{X}$. Then $\mathcal{S}_0 \sqsubseteq \mathcal{C}$.

- Moreover, if $\mathcal{S}_i^* = \mathcal{S}_i$, $1_\Omega \in (\mathcal{S}_i)^\cup$ and $\mathcal{S}_i \sqsubseteq \mathcal{C}$ for $i \geq 0$, then the same holds for $i = i + 1$.
- Thus $\mathcal{S} \sqsubseteq \mathcal{C}$, and hence $\mathcal{S} = \mathcal{C}$.

The main property of the WL-algorithm

Theorem

Let

- $\vec{\mathcal{S}} = (S_1, \dots, S_m)$ and $\vec{\mathcal{T}} = (T_1, \dots, T_m)$ be ordered partitions of Ω^2 and Δ^2 ,
- $\langle\langle \vec{\mathcal{S}} \rangle\rangle = (P_1, \dots, P_k)$ and $\langle\langle \vec{\mathcal{T}} \rangle\rangle = (Q_1, \dots, Q_\ell)$ be the canonical ordering of the coherent closures produced by WL-algorithm,
- $f \in \text{Iso}(\mathcal{S}, \mathcal{T})$ be such that $S_i^f = T_i$, $i = 1, \dots, m$.

Then $k = \ell$ and $P_i^f = Q_i$, $i = 1, \dots, k$. In particular, the mapping $P_i \mapsto Q_i$ is an algebraic isomorphism from $\langle\langle \mathcal{S} \rangle\rangle$ onto $\langle\langle \mathcal{T} \rangle\rangle$.

Thus the ISO can be reformulated as follows: test whether given an algebraic isomorphism is induced by a combinatorial one.

Coherent closure of a Cayley graph

Problem

Given a finite group H and two subsets $S, T \subseteq H$, decide whether $\text{Cay}(H, S) \cong \text{Cay}(H, T)$.

Proposition

The coherent closure $\mathcal{X} = \langle\langle \text{Cay}(H, S) \rangle\rangle$ of the Cayley graph $\text{Cay}(H, S)$ is a fusion scheme of $\text{Inv}(H_R)$.

Proof.

- We have $\text{Aut}(\mathcal{X}) = \text{Aut}(\text{Cay}(H, S)) \geq H_R$.
- Therefore \mathcal{X} is a scheme and $\text{Inv}(\text{Aut}(\mathcal{X})) \sqsubseteq \text{Inv}(H_R)$.
- Thus $\mathcal{X} \sqsubseteq \text{Inv}(\text{Aut}(\mathcal{X})) \sqsubseteq \text{Inv}(H_R)$.

Exercise: the coherent configuration $\text{Inv}(H_R)$ is homogeneous, and any basic relation of it is of the form $h_L \in \text{Sym}(H)$ for some $h \in H$.

Association schemes

Proposition

A pair (Ω, \mathcal{S}) where $\mathcal{S} \vdash \Omega^2$ is an association scheme (=homogeneous coherent configuration) if and only if the following statements hold:

- $1_\Omega \in \mathcal{S}$;
- $\mathcal{S}^* = \mathcal{S}$;
- $\text{bl}(\mathcal{S}) = \mathcal{S}$.

Recall that the intersection numbers c_{RS}^T of the scheme (Ω, \mathcal{S}) are defined so that

$$\forall S, R, T \in \mathcal{S} \quad \forall (\alpha, \beta) \in T : |\alpha S \cap R \beta| = c_{RS}^T.$$

The number $n_S := c_{S S^*}^{1_\Omega} = |\omega S|$ is called the **valency** of S .

Elementary properties of association schemes

Proposition

Let (Ω, \mathcal{S}) be an association scheme. Then

- each relation $R \in \mathcal{S}^\cup$ is regular;
- the set \mathcal{S}^\cup is closed under intersections, unions and complements;
- $1_\Omega, \Omega^2 \in \mathcal{S}^\cup$;
- $(\mathcal{S}^\cup)^* = \mathcal{S}^\cup$;
- \mathcal{S} is closed under relational product;
- for any $S \in \mathcal{S}$ there exists m such that $1_\Omega \in S^m$;
- if $S \in \mathcal{S}^\cup$ and $S^+ = \bigcup_{i=0}^\infty S^i$, then $S^+ \in \mathcal{S}^\cup$;
- if $S \in \mathcal{S}^\cup$, then S^+ is an equivalence relation on Ω .

Lecture 5.

Main classes of association schemes

Definition

A scheme is $\mathcal{X} = (\Omega, \mathcal{S})$ is called

- **symmetric** if every $S \in \mathcal{S}$ is symmetric,
- **antisymmetric** if every $S \in \mathcal{S}$, $S \neq 1_\Omega$, is anti-symmetric,
- **commutative** if the algebra $\mathbb{F}\mathcal{X}$ is commutative.

Exercise: a symmetric scheme is always commutative.

Proposition

Let (Ω, \mathcal{S}) be a scheme. Then

- $\sum_{S \in \mathcal{S}} n_S = |\Omega|$;
- $1 \leq |\mathcal{S}| \leq |\Omega|$;
- $|\mathcal{S}| = 2$ if and only if $\mathcal{S} = \{1_\Omega, \Omega^2 \setminus 1_\Omega\}$;
- $|\mathcal{S}| = |\Omega|$ if and only if $n_S = 1$ for all $S \in \mathcal{S}$.

Primitive and imprimitive schemes

Definition

A scheme $\mathcal{X} = (\Omega, \mathcal{S})$ is called **imprimitive** if \mathcal{S}^\cup contains an equivalence relation E other than 1_Ω and Ω^2 ; the sets in Ω/E form a partition of Ω called the **imprimitivity system** of \mathcal{S} .

Proposition

Given an equivalence relation $E \in \mathcal{S}^\cup$ we have $|\Omega/E| \cdot n_E = |\Omega|$.

Proposition

The following statements are equivalent:

- \mathcal{X} is imprimitive;
- there exists $S \in \mathcal{S}$ such that $1_\Omega \neq S^+ \neq \Omega^2$;
- there exists $\mathcal{T} \subset \mathcal{S}$ with $1 < |\mathcal{T}| < |\mathcal{S}|$ such that $\langle A(T) \rangle_{T \in \mathcal{T}}$ is a subalgebra of $\mathbb{F}\mathcal{X}$.

Schurian schemes

Recall that a scheme $\mathcal{X} = (\Omega, \mathcal{S})$ is **schurian** if there exists a group $G \leq \text{Sym}(\Omega)$ such that $\mathcal{X} = \text{Inv}(G)$. In this case

- G is transitive on Ω ;
- $\omega S \in \text{Orb}(G_\omega, \Omega)$ for all $\omega \in \Omega$ and $S \in \mathcal{S}$;
- the mapping $S \mapsto \{g \in G : \omega^g \in \omega S\}$ is a bijection from \mathcal{S} onto the set of double cosets of G_ω in G ;
- a rescaling of the above mapping is an isomorphism from $\mathbb{F}\mathcal{X}$ onto the Hecke algebra $\mathbb{F}(G_\omega \backslash G / G_\omega)$;
- G is primitive if and only if $\text{Inv}(G)$ is primitive.

Thin schemes

Definition

A scheme (Ω, \mathcal{S}) is called **thin** if $n_S = 1$ for each $S \in \mathcal{S}$.

Proposition

A scheme (Ω, \mathcal{S}) is thin if and only if \mathcal{S} is a regular subgroup of $\text{Sym}(\Omega)$.

Example:

- (H, H_R) and (H, H_L) are thin schemes;
- the bijection $h \mapsto h^{-1}$ is an isomorphism from (H, H_R) onto (H, H_L) ;
- $\text{Inv}(H_R) = H_L$ and $\text{Inv}(H_L) = H_R$;
- $\text{Aut}(H_L) = H_R$ and $\text{Aut}(H_R) = H_L$.

Cayley schemes

Definition

An association scheme \mathcal{X} which is a fusion of (H, H_L) is called a **Cayley scheme** over H , or equivalently, $\text{Aut}(\mathcal{X}) \geq H_R$.

A basic relation S of such a scheme is a Cayley graph $\text{Cay}(H, Se)$.

Proposition

Let (H, \mathcal{R}) be a Cayley scheme. Then the set $\mathcal{S} := \{Re : R \in \mathcal{R}\}$ is a partition of H , and

- $\{e\} \in \mathcal{S}$;
- $X^{(-1)} \in \mathcal{S}$ for all $X \in \mathcal{S}$;
- for any triple $X, Y, Z \in \mathcal{S}$ the number $c_{XY}^Z := |Y \cap X^{(-1)}z|$ does not depend on $z \in Z$.

Cayley schemes and Schur partitions

Definition

A partition \mathcal{S} of a group H is called a **Schur partition** if it satisfies the above conditions, that is

- $\{e\} \in \mathcal{S}$;
- $X^{(-1)} \in \mathcal{S}$ for all $X \in \mathcal{S}$;
- for any triple $X, Y, Z \in \mathcal{S}$ the number $c_{XY}^Z := |Y \cap X^{(-1)}z|$ does not depend on $z \in Z$.

Notice that $|Y \cap X^{(-1)}z| = \{(y, x) \in Y \times X : yx = z\}$.

Proposition

Let \mathcal{S} be a partition of H . Then the partition $\text{Cay}(H, \mathcal{S}) := \{\text{Cay}(H, X) : X \in \mathcal{S}\}$ of the set $H \times H$ is a scheme if and only if \mathcal{S} is a Schur partition.

Group rings

Notation

- RH is the group algebra of H over a unitary ring R ;
- if $x = \sum_h x_h h$ and $y = \sum_h y_h h$ belong to RH , then

$$xy := \sum_{h,f \in H} x_h y_f (hf) \quad \text{and} \quad x \circ y = \sum_{h \in H} (x_h y_h) h$$

are the **convolution** and the **Schur-Hadamard** product of x and y ;

- any \circ -idempotent is of the form $\underline{X} := \sum_{h \in X} h$ where $X \subseteq H$;
- $\{\underline{h}\}$ is abbreviated as h ;
- if $\mathcal{S} \vdash H$, then $\underline{\mathcal{S}} := \{\underline{X} : X \in \mathcal{S}\}$;
- for each $m \in \mathbb{Z}$ and $x \in RH$ we set $x^{(m)} := \sum_{h \in H} x_h h^m$.

Schur partitions and Schur rings

Proposition

Let $\mathcal{S} \vdash H$ be such that $\{e\} \in \mathcal{S}$ and $\mathcal{S}^{(-1)} = \mathcal{S}$. Then \mathcal{S} is a Schur partition if and only if the linear span $\langle \underline{\mathcal{S}} \rangle$ is a subalgebra of $\mathbb{Q}H$.

Definition

A subalgebra $\mathcal{A} \leq \mathbb{Q}H$ is called a **Schur ring** over H if there exists a Schur partition $\mathcal{S} \vdash H$ such that $\mathcal{A} = \langle \underline{\mathcal{S}} \rangle$. The elements of \mathcal{S} and \mathcal{S}^{\cup} are called **basic sets** and **\mathcal{A} -sets** respectively.

Theorem

A vector space $\mathcal{A} \leq \mathbb{Q}H$ is a Schur ring if and only if $e, \underline{H} \in \mathcal{A}$ and \mathcal{A} closed with respect to the convolution, \circ and (-1) .

Generating S-ring

Proposition

If $\mathcal{A} = \langle \underline{\mathcal{S}} \rangle$ and $\mathcal{B} = \langle \underline{\mathcal{T}} \rangle$ are S-rings, then

- $\mathcal{A} \subseteq \mathcal{B}$ if and only if $\mathcal{S} \subseteq \mathcal{T}$;
- $\mathcal{A} \cap \mathcal{B} = \langle \underline{\mathcal{S} \wedge \mathcal{T}} \rangle$.

The intersection of all S-rings containing elements $x, y, z, \dots \in \mathbb{Q}H$ is denoted by $\langle\langle x, y, z, \dots \rangle\rangle$.

Theorem

Let $X \subseteq H$, $\mathcal{A} = \langle\langle \underline{X} \rangle\rangle$ a Schur ring generated by \underline{X} and \mathcal{S} the corresponding S-partition. Then $\langle\langle \text{Cay}(H, X) \rangle\rangle = \text{Cay}(H, \mathcal{S})$ and $S \in \mathcal{S}^{\cup}$.

Schur-Wielandt principle

For a function $f : \mathbb{Q} \rightarrow \mathbb{Q}$ and $x = \sum_{h \in H} x_h h$ set

$$f[x] := \sum_{h \in H} f(x_h) h.$$

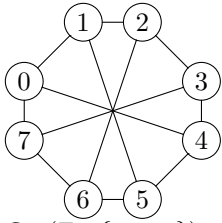
Proposition

Let $f : \mathbb{Q} \rightarrow \mathbb{Q}$ be a function, \mathcal{A} is an S-ring over H and $x \in \mathcal{A}$. Then $f[x] \in \mathcal{A}$.

Corollary

Let \mathcal{A} be an S-ring over H , $x \in \mathcal{A}$ and $r \in \mathbb{Q}$. Then the element \underline{X} where $X = \{h \in H : x_h = r\}$, belongs to \mathcal{A} . In particular, X is an \mathcal{A} -set.

Proof. Follows from the Proposition because $X = f[x]$ where $f = \delta_r$ is the Kronecker delta-function.



Cay($\mathbb{Z}_8, \{1, 4, 7\}$)

Computation of $\langle\langle \underline{S} \rangle\rangle$ where $S = \{1, 4, 7\} \subseteq \mathbb{Z}_8$

Write $\underline{S} = \{1, 4, 7\}$ as $c + c^4 + c^7$ where $c^8 = 1$. Then

- $\langle\langle \underline{S} \rangle\rangle$ contains $\underline{S}^2 = 3c^0 + c^2 + c^6 + 2c^5 + 2c^3$.
- Therefore $c^5 + c^3, c^2 + c^6 \in \langle\langle \underline{S} \rangle\rangle$.
- So $\langle\langle \underline{S} \rangle\rangle$ contains $(c^2 + c^6)^2 = 2c^0 + 2c^4$.
- Thus $\langle\langle \underline{S} \rangle\rangle = \langle c^0, c^4, c + c^7, c^2 + c^6, c^3 + c^5 \rangle$, and
- $\langle\langle \text{Cay}(\mathbb{Z}_8, \{1, 4, 7\}) \rangle\rangle = \text{Cay}(\mathbb{Z}_8, \{\{0\}, \{4\}, \{1, 7\}, \{2, 6\}, \{3, 5\}\})$.

In other words, we found the coherent closure of the graph:

Examples: Schur partitions over the group \mathbb{Z}_8

The following list was generated by the computer program COCO (thanks to Misha Klin).

- $\{0\}, \{1, 2, 3, 4, 5, 6, 7\};$
- $\{0\}, \{1, 3, 5, 7\}, \{2, 6, 4\};$
- $\{0\}, \{1, 3, 5, 7, 2, 6\}, \{4\};$
- $\{0\}, \{1, 3, 5, 7\}, \{2, 6\}, \{4\};$
- $\{0\}, \{1, 3, 5, 7\}, \{2\}, \{6\}, \{4\};$
- $\{0\}, \{1, 5\}, \{3, 7\}, \{2\}, \{6\}, \{4\};$
- $\{0\}, \{1, 5\}, \{3, 7\}, \{2, 6\}, \{4\};$
- $\{0\}, \{1, 3\}, \{5, 7\}, \{2, 6\}, \{4\};$
- $\{0\}, \{1, 7\}, \{3, 5\}, \{2, 6\}, \{4\};$
- $\{0\}, \{1\}, \{2\}, \{3\}, \{4\}, \{5\}, \{6\}, \{7\};$

Further examples

Proposition

Let $F \leq \text{Aut}(H) \leq \text{Sym}(H)$. Then the orbit partition $\text{Orb}(F, H)$ is a Schur partition. The corresponding S-ring coincides with $\mathbb{Q}[H]^F$.

Partial cases

- $F = \text{Inn}(H) \implies \mathbf{Z}(\mathbb{Q}[H])$ is an S-ring. Its basic sets coincide with conjugacy classes of H . Fusion S-rings of $\mathbf{Z}(\mathbb{Q}[H])$ are in one-to-one correspondence with supercharacters introduced recently by Isaacs et. el.
- let R be a ring, $H = (R, +)$ and $K \leq R^\times$. The corresponding S-ring $\mathbb{Q}[H]^K$ is called **cyclotomic**. Its basic sets have a form $Kr, r \in R$.

Further examples

Proposition (subgroup S-rings)

\mathcal{L} be a sublattice of a subgroup lattice of H which contains $\{e\}$ and H . If any two subgroup $K, L \in \mathcal{L}$ are permutable, then $\langle \underline{L} \rangle_{L \in \mathcal{L}}$ is a Schur ring.

Hecke algebras

Let $K \leq H$ be an arbitrary subgroup and $\mathcal{S} = \{KhK \mid h \in H\}$ be a partition of H into double cosets of K . The linear span $\underline{\mathcal{S}}$ is known as **Hecke algebra** w.r.t. K . It is closed w.r.t. $(^{-1}), \circ, \cdot$ but doesn't contain e .

Properties of S-rings

Proposition

Let \mathcal{S} be a Schur partition of H and $\text{Cay}(H, \mathcal{S})$ the corresponding Cayley scheme. Then $\text{Cay}(H, \mathcal{S})^\cup = \text{Cay}(H, \mathcal{S}^\cup)$ and

- the set \mathcal{S}^\cup is closed w.r.t. boolean operations;
- $\{e\}, H \in \mathcal{S}^\cup$;
- $(\mathcal{S}^\cup)^* = \mathcal{S}^\cup$;
- \mathcal{S} is closed w.r.t. group product;
- $S \in \mathcal{S}^\cup \implies \langle S \rangle \in \mathcal{S}^\cup$;

A relation $E = \text{Cay}(H, \mathcal{S}), S \in \mathcal{S}^\cup$ is an equivalence iff S is a subgroup of H .

Definition

A subgroup $F \leq H$ is called an \mathcal{A} -subgroup if $F \in \mathcal{A}$. An S-ring is called **primitive** iff $\{e\}, H$ are the only \mathcal{A} -subgroups.

Schurian S-rings

Theorem (Schur)

Let H be a group and $H_R \leq G \leq \text{Sym}(H)$. Then the orbits of G_e form an S-partition.

Proof. Set $\mathcal{S} := \text{Inv}(G)$. Then $H_R \leq G \implies \mathcal{S} \sqsubseteq \text{Inv}(H_R) = H_L$. Thus \mathcal{S} is a Cayley scheme. Hence $e\mathcal{S} = \{eS \mid S \in \mathcal{S}\}$ is a Schur partition of H . Since \mathcal{S} is schurian, $e\mathcal{S} = \text{Orb}(G_e, H)$. \square

An S-partition is called **Schurian** if it has a form $\text{Orb}(G_e, H)$ for some $G, H_R \leq G \leq \text{Sym}(H)$

Subgroup factorization and Schur rings

Theorem (Schur)

Let $G = AH$ be a factorization into a subgroup product with $A \cap H = \{e\}$. Then the subalgebra

$$\mathbf{C}_{\mathbb{Q}[H]}(\underline{A}) := \{x \in \mathbb{Q}[H] \mid x\underline{A} = \underline{A}x\}.$$

is a Schur ring the basic sets of which have the form $AhA \cap H$.

A concrete example

A simple group $PSL_3(2)$ has a decomposition into a product AH where $A \cong D_8$ and $H \cong F_{21}$. The corresponding S-ring over H has rank six and its Cayley scheme is isomorphic to a flag scheme of a projective plane of order 2.

Isomorphisms between Schur rings

Let $\mathcal{S} \vdash H$ and $\mathcal{T} \vdash K$ be two S-partitions of groups H and K resp. The S-rings $\mathcal{A} := \langle \mathcal{S} \rangle, \mathcal{B} := \langle \mathcal{T} \rangle$ are

- **Cayley isomorphic**, notation \cong_{Cay} , if there exists a group isomorphism $f : H \rightarrow K$ s.t. $\mathcal{S}^f = \mathcal{T}$;
- **combinatorially isomorphic** if the schemes $\text{Cay}(H, \mathcal{S})$ and $\text{Cay}(K, \mathcal{T})$ are isomorphic (as schemes);
- **algebraically isomorphic** if the schemes $\text{Cay}(H, \mathcal{S})$ and $\text{Cay}(K, \mathcal{T})$ are algebraically isomorphic.

In what follows we abbreviate

$$\text{Aut}(\mathcal{A}) := \text{Aut}(\text{Cay}(H, \mathcal{S})), \text{Iso}(\mathcal{A}) := \text{Iso}(\text{Cay}(H, \mathcal{S})).$$

Isomorphisms between S-rings

Proposition

$\mathcal{A} \cong_{\text{alg}} \mathcal{B}$ iff there exists a bijection $f : \mathcal{S} \rightarrow \mathcal{T}$ s.t. $c_{PQ}^R = c_{P^f Q^f}^{R^f}$.

Proposition

$\mathcal{A} \cong_{\text{com}} \mathcal{B}$ iff there exists a bijection $f : H \rightarrow K$ s.t. $(e_H)^f = e_K$ and

- for any $h \in H$ and $S \in \mathcal{S}$ it holds that $(hS)^f = h^f S^f$;
- $\mathcal{S}^f = \mathcal{T}$;
- $f|_{\mathcal{S}}$ is an algebraic isomorphism between \mathcal{A} and \mathcal{B}

$$\begin{aligned} \mathcal{A} \cong_{\text{Cay}} \mathcal{B} &\implies \mathcal{A} \cong \mathcal{B} \implies \mathcal{A} \cong_{\text{alg}} \mathcal{B}. \\ \mathcal{A} \cong_{\text{Cay}} \mathcal{B} \neq \mathcal{A} \cong \mathcal{B} &\neq \mathcal{A} \cong_{\text{alg}} \mathcal{B}. \end{aligned}$$

Application to Cayley graph isomorphism problem.

- Let $f : \text{Cay}(H, S) \rightarrow \text{Cay}(H, T)$ be an isomorphism s.t. $f(e) = e$;
- then $\mathcal{S}^f = \mathcal{T}$ where \mathcal{S} and \mathcal{T} are S-partitions generated by \underline{S} and \underline{T} resp.;
- f^* is an algebraic isomorphism between S-rings $\underline{\mathcal{S}}$ and $\underline{\mathcal{T}}$.

Klin-Pöschel approach.

How to solve ISO for Cayley graphs over a finite group H .

- Find all S-rings over H . Let $\mathcal{A}_1, \dots, \mathcal{A}_N$ be the complete list of them;
- For all pairs i, j find the set Φ_{ij} of algebraic isomorphisms between them;
- For each $\phi \in \Phi_{ij}$ find a combinatorial isomorphism f between \mathcal{A}_i and \mathcal{A}_j s.t. $f^* = \phi$ (if such f exists);
- Collect all permutations f found on the previous stage. Let P be the set of all those permutations.

Proposition

The set P constructed above is a [solving set](#) for the Cayley graphs over H , that is two Cayley graphs $\text{Cay}(H, S)$ and $\text{Cay}(H, T)$ are isomorphic iff there exists $f \in P$ s.t. $\text{Cay}(H, S)^f = \text{Cay}(H, T)$.

Example

N	S-partition \mathcal{S}	$ \text{Alg}(\mathcal{S}) $	$\text{Iso}(\mathcal{S})/\text{Aut}(\mathcal{S})$ transversal
1	$\{0\}, \{1, 2, 3, 4, 5, 6, 7\}$	1	μ_1
2	$\{0\}, \{1, 3, 5, 7\}, \{2, 6, 4\}$	1	μ_1
3	$\{0\}, \{1, 3, 5, 7, 2, 6\}, \{4\}$	1	μ_1
4	$\{0\}, \{1, 3, 5, 7\}, \{2, 6\}, \{4\}$	1	μ_1
5	$\{0\}, \{1, 3, 5, 7\}, \{2\}, \{6\}, \{4\}$	2	μ_1, μ_3
6	$\{0\}, \{1, 5\}, \{3, 7\}, \{2\}, \{6\}, \{4\}$	4	$\mu_1, \mu_3, \sigma, \sigma\mu_3$
7	$\{0\}, \{1, 5\}, \{3, 7\}, \{2, 6\}, \{4\}$	2	μ_1, μ_3
8	$\{0\}, \{1, 3\}, \{5, 7\}, \{2, 6\}, \{4\}$	2	μ_1, μ_5
9	$\{0\}, \{1, 7\}, \{3, 5\}, \{2, 6\}, \{4\}$	2	μ_1, μ_3
10	$\{0\}, \{1\}, \{2\}, \{3\}, \{4\}, \{5\}, \{6\}, \{7\}$	4	$\mu_1, \mu_3, \mu_5, \mu_7$

Here $\sigma = (2, 6)(3, 7)$ and μ_a is an automorphism of $\mathbb{Z}_8: x \mapsto ax$. Thus $\{\mu_1, \mu_3, \mu_5, \mu_7, \sigma, \sigma\mu_3\}$ is a solving set for \mathbb{Z}_8 .

Solving sets for cyclic groups

Theorem

Two S-rings over cyclic groups are algebraically isomorphic iff they coincide.

This implies the following modification of the original Klin-Pöschel approach

- Find all S-rings over \mathbb{Z}_n , let $\mathcal{A}_1, \dots, \mathcal{A}_N$ be the complete list of them;
- For each \mathcal{A}_i find a transversal T_i of $\text{Iso}(\mathcal{A}_i)/\text{Aut}(\mathcal{A}_i)$
- Then the union of T_i produces a solving set for \mathbb{Z}_n .

Theorem

Given a number n , one can construct a solving set for \mathbb{Z}_n of at most n^3 permutations in time $n^{O(1)}$.

Lecture 6.

Cyclotomic S-rings

Proposition

Let $F \leq \text{Aut}(H) \leq \text{Sym}(H)$. Then the orbit partition $\text{Orb}(F, H)$ is a Schur partition; the corresponding S-ring coincides with $(\mathbb{Q}H)^F$.

Special cases

- If $F = \text{Inn}(H)$, then $(\mathbb{Q}H)^F = \mathbf{Z}(\mathbb{Q}H)$ and the basic sets are the conjugacy classes of H . Fusion S-rings of $\mathbf{Z}(\mathbb{Q}H)$ are in one-to-one correspondence with [supercharacters](#).
- If R is a finite ring, $H = (R, +)$ and $K \leq R^\times$, then $\text{Cyc}(K, H) := (\mathbb{Q}H)^K$ is called a [cyclotomic](#) S-ring; its basic sets are rK , $r \in R$.
- When $R = \mathbb{Z}_n$, then H is a cyclic group and K isomorphic to a subgroup of $\text{Aut}(H)$; the S-ring $\text{Cyc}(K, H)$ is called [circulant](#).

Properties of \mathcal{A} -subsets

Proposition

Let \mathcal{S} be a Schur partition of H and $\text{Cay}(H, \mathcal{S})$ the corresponding Cayley scheme. Then $\text{Cay}(H, \mathcal{S})^\cup = \text{Cay}(H, \mathcal{S}^\cup)$ and

- \mathcal{S}^\cup is closed with respect to boolean operations;
- $\{e\}, H \in \mathcal{S}^\cup$;
- $(\mathcal{S}^\cup)^{-1} = \mathcal{S}^\cup$;
- \mathcal{S}^\cup is closed with respect to the group product;
- $\langle X \rangle \in \mathcal{S}^\cup$ for all $X \in \mathcal{S}^\cup$;
- if $X \in \mathcal{S}^\cup$, then $E = \text{Cay}(H, X)$ is an equivalence relation if and only if $X \leq H$.

Definition

A subgroup $F \leq H$ is called an \mathcal{A} -group if $F \in \mathcal{A}$; \mathcal{A} is called **primitive** if $\{e\}, H$ are the only \mathcal{A} -subgroups.

Schurian S-rings

Theorem (Schur, 1933)

Let H be a group and $H_R \leq G \leq \text{Sym}(H)$. Then $\text{Orb}(G_e, H)$ is an S-partition.

Proof.

- Set $\mathcal{X} := \text{Inv}(G)$.
- Since $H_R \leq G$, we have $\mathcal{X} \sqsubseteq \text{Inv}(H_R) = H_L$.
- Thus $\mathcal{X} = (H, \mathcal{R})$ is a Cayley scheme over H ,
- and $e\mathcal{R} = \{eS : S \in \mathcal{R}\}$ is a Schur partition of H .
- Since \mathcal{X} is schurian, $e\mathcal{R} = \text{Orb}(G_e, H)$.

Definition

An S-partition is called **Schurian** if it has a form $\text{Orb}(G_e, H)$ for some G such that $H_R \leq G \leq \text{Sym}(H)$.

Isomorphisms between S-rings: definitions

Let $\mathcal{S} \vdash H$ and $\mathcal{T} \vdash K$ be S-partitions of groups H and K .

Definition

The S-rings $\mathcal{A} := \langle \mathcal{S} \rangle$ and $\mathcal{B} := \langle \mathcal{T} \rangle$ are called

- **Cayley isomorphic**, notation \cong_{Cay} , if there exists a group isomorphism $f : H \rightarrow K$ such that $\mathcal{S}^f = \mathcal{T}$;
- **combinatorially isomorphic** if the schemes $\text{Cay}(H, \mathcal{S})$ and $\text{Cay}(K, \mathcal{T})$ are isomorphic (as schemes);
- **algebraically isomorphic** if the schemes $\text{Cay}(H, \mathcal{S})$ and $\text{Cay}(K, \mathcal{T})$ are algebraically isomorphic.

In what follows we abbreviate

$$\text{Aut}(\mathcal{A}) := \text{Aut}(\text{Cay}(H, \mathcal{S})), \quad \text{Iso}(\mathcal{A}) := \text{Iso}(\text{Cay}(H, \mathcal{S})).$$

Isomorphisms between S-rings: characterization

Proposition

Let $\mathcal{A} := \langle \mathcal{S} \rangle$ and $\mathcal{B} := \langle \mathcal{T} \rangle$ be S-rings. Then

- $\mathcal{A} \cong_{\text{alg}} \mathcal{B}$ if and only if there exists a bijection $f : \mathcal{S} \rightarrow \mathcal{T}$ such that $c_{PQ}^R = c_{P^f Q^f}^{R^f}$ for all $P, Q, R \in \mathcal{S}$.
- $\mathcal{A} \cong_{\text{com}} \mathcal{B}$ if and only if there exists a bijection $f : H \rightarrow K$ such that $(e_H)^f = e_K$, $f|_{\mathcal{S}}$ is an algebraic isomorphism from \mathcal{A} onto \mathcal{B} and $(hS)^f = h^f S^f$ for all $h \in H$ and $S \in \mathcal{S}$.

Remark:

$$\begin{aligned} \mathcal{A} \cong_{\text{Cay}} \mathcal{B} &\Rightarrow \mathcal{A} \cong \mathcal{B} \Rightarrow \mathcal{A} \cong_{\text{alg}} \mathcal{B}, \\ \mathcal{A} \cong_{\text{Cay}} \mathcal{B} &\not\Rightarrow \mathcal{A} \cong \mathcal{B} \not\Rightarrow \mathcal{A} \cong_{\text{alg}} \mathcal{B}. \end{aligned}$$

Klin-Pöschel approach

A “solution” to the ISO for Cayley graphs over a finite group H :

- find all S-rings over H ; let $\mathcal{A}_1, \dots, \mathcal{A}_N$ be the complete list of them;
- for all pairs i, j find the set Φ_{ij} of algebraic isomorphisms between them;
- for each $\phi \in \Phi_{ij}$ find a combinatorial isomorphism f between \mathcal{A}_i and \mathcal{A}_j such that $f^* = \phi$ (if such f exists);
- collect all permutations f found on the previous stage; let P be the set of all those permutations.

Proposition

The set P is a [solving set](#) for the Cayley graphs over H : two Cayley graphs $\text{Cay}(H, S)$ and $\text{Cay}(H, T)$ are isomorphic if and only if there exists $f \in P$ such that $\text{Cay}(H, S)^f = \text{Cay}(H, T)$.

Example

N	S-partition \mathcal{S}	$ \text{Alg}(\mathcal{S}) $	$\text{Iso}(\mathcal{S})/\text{Aut}(\mathcal{S})$ transversal
1	$\{0\}, \{1, 2, 3, 4, 5, 6, 7\}$	1	μ_1
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4	$\{0\}, \{1, 3, 5, 7\}, \{2, 6\}, \{4\}$	1	μ_1
5	$\{0\}, \{1, 3, 5, 7\}, \{2\}, \{6\}, \{4\}$	2	μ_1, μ_3
6	$\{0\}, \{1, 5\}, \{3, 7\}, \{2\}, \{6\}, \{4\}$	4	$\mu_1, \mu_3, \sigma, \sigma\mu_3$
7	$\{0\}, \{1, 5\}, \{3, 7\}, \{2, 6\}, \{4\}$	2	μ_1, μ_3
8	$\{0\}, \{1, 3\}, \{5, 7\}, \{2, 6\}, \{4\}$	2	μ_1, μ_5
9	$\{0\}, \{1, 7\}, \{3, 5\}, \{2, 6\}, \{4\}$	2	μ_1, μ_3
10	$\{0\}, \{1\}, \{2\}, \{3\}, \{4\}, \{5\}, \{6\}, \{7\}$	4	$\mu_1, \mu_3, \mu_5, \mu_7$

Here $\sigma = (2, 6)(3, 7)$ and μ_a is an automorphism of \mathbb{Z}_8 : $x \mapsto ax$. Thus $\{\mu_1, \mu_3, \mu_5, \mu_7, \sigma, \sigma\mu_3\}$ is a solving set for \mathbb{Z}_8 .

Solving sets for cyclic groups

Theorem (Muzychuk, 1994)

Two S-rings over cyclic groups are algebraically isomorphic if and only if they coincide.

This gives the following modification of the original Klin-Pöschel approach:

- find all S-rings over \mathbb{Z}_n , let $\mathcal{A}_1, \dots, \mathcal{A}_N$ be the complete list of them;
- for each \mathcal{A}_i find a transversal T_i of $\text{Iso}(\mathcal{A}_i)/\text{Aut}(\mathcal{A}_i)$;
- now the union of T_i 's produces a solving set for \mathbb{Z}_n .

Theorem (Muzychuk, 2004)

Given a number n , one can construct a solving set for \mathbb{Z}_n of at most n^3 permutations in time $n^{O(1)}$.

B-groups

Definition (Wielandt, 1964)

A finite group H is called a [B-group](#) if every primitive permutation group G satisfying $H_R \leq G \leq \text{Sym}(H)$ is 2-transitive.

Theorem

A group H is a B-group if and only if every primitive Schurian S-ring over H has rank 2.

To prove this theorem, it suffices to verify the following statement.

Proposition

Let $H_R \leq G \leq \text{Sym}(H)$, $\mathcal{A} := \langle \text{Orb}(G_e, H) \rangle$ the corresponding S-ring and $B \subseteq H$ with $e \in B$. Then B is a block of G if and only if B is an \mathcal{A} -group.

Proof of the theorem on B-groups

Proof.

- Assume first that B is a block of G .
- Then B is a block of H_R .
- So for any $h \in H$: $Bh = B$ or $Bh \cap B = \emptyset$.
- It follows that B is a coset.
- Since $e \in B$, this implies that $B \leq H$.
- Besides, since B is a block, $B^{G_e} = B$.
- Thus B is a union of G_e -orbits, and hence is an \mathcal{A} -group.
- Conversely, let now B be an \mathcal{A} -group.
- Then $\text{Cay}(H, B)$ is a G -invariant equivalence relation.
- Therefore the classes of $\text{Cay}(H, B)$ are blocks of G ,
- and B is a block of G .

Exercise:

prove that a simple group of order > 2 is not a B-group.

Strong B-groups

Definition

A finite group H is **strong B-group** if the trivial S-ring is the only primitive S-ring over H .

Known strong B-groups:

- a cyclic group of composite order (Schur, 1933);
- an abelian group with a cyclic Sylow subgroup (Wielandt, 1935);
- an abelian group with a Sylow p -subgroup of type (p^a, p^b) , $a < b$ (Bercov, 1962);
- a dihedral group (Wielandt, 1949).

Exercise:

prove that any group of order 6 or 8 is a strong B-group.

Lecture 7.

\mathcal{A} -groups

Let $\mathcal{A} \leq \mathbb{Q}H$ be an S-ring over H and \mathcal{S} the corresponding S-partition of H .

Proposition

Given \mathcal{A} -groups A and B , $A \cap B$ and $\langle A, B \rangle$ are \mathcal{A} -groups.

Proposition.

Given $S \in \mathcal{S}^\cup$, the groups $\langle S \rangle$ and

$$\begin{aligned}\text{Rad}_\ell(S) &:= \{h \in H \mid hS = S\}, \\ \text{Rad}_r(S) &:= \{h \in H \mid Sh = S\}, \\ \text{Rad}(S) &:= \text{Rad}_\ell(S) \cap \text{Rad}_r(S)\end{aligned}$$

\mathcal{A} -groups.

Indeed, $h \in \text{Rad}_\ell(S)$ if and only if $Sh = S$ if and only if h appears $|S|$ times in the product $\underline{S} \cdot \underline{S}^{(-1)}$. Thus $\text{Rad}_\ell(S) = \delta_{|S|}[\underline{S} \cdot \underline{S}^{(-1)}]$ and by Schur-Wielandt principle $\text{Rad}_\ell(S) \in \mathcal{A}$.

Subring of an S-ring

In what follows $K \leq H$ is an \mathcal{A} -group.

Proposition

The vector space $\mathcal{A}_K := \mathcal{A} \cap \mathbb{Q}K$ is an S-ring over K the basic sets of which are $\mathcal{S}_K := \{S \in \mathcal{S} : S \subseteq K\}$.

The quotient S-ring: auxiliary lemma

Proposition

For basic sets $S, T \in \mathcal{S}$ the following conditions hold:

- (a) $SK \neq TK \implies SK \cap TK = \emptyset$;
- (b) $KSK \neq KTK \implies KSK \cap KTK = \emptyset$;
- (c) $\underline{K} \cdot \underline{S} = c\underline{KS}$ where $c = |Ks \cap S|, s \in S$
- (d) $\underline{K} \cdot \underline{S} \cdot \underline{K} = c\underline{KSK}$ where $c = |KS \cap sK|, s \in S$.

Proof: (a) if $SK \cap TK \neq \emptyset$, then $sk_1 = tk_2$ for some $s \in S, t \in T$ and $k_1, k_2 \in K$. Then $S \cap TK \neq \emptyset$, and $S \subseteq TK \implies SK \subseteq TK$. By symmetry $TK \subseteq SK$.

(c) The coefficient of h in the product $\underline{K} \cdot \underline{S}$ is equal to the intersection number $|K^{(-1)}h \cap S| = |Kh \cap S|$. Since S is a basic set, the number $c := |K^{(-1)}s \cap S|$ doesn't depend on a choice $s \in S$. Thus if $h \in KS$, then $h = ks$ for some $k \in K$ and $s \in S$, and $|Kh \cap S| = |Ks \cap S| = c$.

The quotient S-ring: definition

Denote by $K \backslash \mathcal{S} / K$ the partition of H into classes $KSK, S \in \mathcal{S}$.

Proposition

The linear span $\langle \underline{X} : X \in K \backslash \mathcal{S} / K \rangle$ is closed under \circ, \cdot and (-1) and contains \underline{H} . It coincides with $e_K \mathcal{A} e_K$ where $e_K = \frac{1}{|K|} \underline{K}$.

If K is normal in H then the algebra $e_K \mathcal{A} e_K$ may be identified with an S-ring over H/K . The mapping $\pi : x \mapsto e_K x, x \in \mathbb{Q}H$ is an epimorphism from $\mathbb{Q}H$ onto $e_K \mathbb{Q}H$. Identifying the elements he_K with the cosets hK we obtain the following statement.

Proposition

The image $\pi(\mathcal{A})$ is an S-ring over H/K the basic sets of which are $\pi(T), T \in \mathcal{S}$. This S-ring is called a **quotient** of \mathcal{A} and is denoted as \mathcal{A}/K or $\mathcal{A}_{H/K}$.

Exercises

Exercise 1:

let $C = \langle c \rangle$ be a cyclic group of order 8. Then the subspace $\mathcal{A} = \langle c^0, c^2, c^4, c^6, c + c^5, c^3 + c^7 \rangle$ is an S-ring over C . Check that

- every subgroup of C is an \mathcal{A} -group;
- $\mathcal{A}_{\langle c^2 \rangle} = \mathbb{Q}[c^2]$;
- $\mathcal{A}/\langle c^4 \rangle = \mathbb{Q}[C/\langle c^4 \rangle]$;
- find $\text{Aut}(\mathcal{A}), \text{Iso}(\mathcal{A})$ and $\text{Alg}(\mathcal{A})$.

Exercise 2:

let $\mathcal{A} \leq \mathbb{Q}H$ be an S-ring and $K \trianglelefteq H$ an \mathcal{A} -subgroup. Then \mathcal{A}_K and $\mathcal{A}_{H/K}$ are schurian, whenever so is \mathcal{A} .

Wreath product of S-rings

Proposition

$$\dim(\mathcal{A}) \geq \dim(\mathcal{A}_K) + \dim(\mathcal{A}_{H/K}) - 1.$$

Proof. \mathcal{A}_K and $e_K \mathcal{A}$ are subspaces of \mathcal{A} . Since $\dim(\mathcal{A}_{H/K}) = \dim(e_K \mathcal{A})$ and $\mathcal{A}_K \cap e_K \mathcal{A} = \langle \underline{K} \rangle$, we are done.

Proposition

The following conditions are equivalent:

- (a) $\dim(\mathcal{A}) = \dim(\mathcal{A}_K) + \dim(\mathcal{A}_{H/K}) - 1$;
- (b) $\mathcal{A} = \mathcal{A}_K + e_K \mathcal{A}$;
- (c) $K \leq \text{Rad}(T)$ for each basic set of \mathcal{A} with $T \cap K = \emptyset$.

If any of the above conditions hold, the S-ring \mathcal{A} is called the **wreath product** of \mathcal{A}_K by $\mathcal{A}_{H/K}$, notation $\mathcal{A}_K \wr \mathcal{A}_{H/K}$.

Generalized wreath product

Definition

Let $1 \leq L \leq U \leq H$ and $L \trianglelefteq H$ be \mathcal{A} -groups. The S-ring \mathcal{A} is called a **generalized wreath product** or **U/L -wreath product** of \mathcal{A}_U and $\mathcal{A}_{U/L}$ if every basic set $S \in \mathcal{S}$ outside of U is a union of L -cosets. The product is called **nontrivial** or **proper** if $L \neq 1, U \neq H$.

Remark: under above conditions $(\mathcal{A}_U)_{U/L} = \mathcal{A}_{U/L} = (\mathcal{A}/L)_{U/L}$.

Theorem

Let $S = U/L$ be a section of a group H , and let \mathcal{A}_1 and \mathcal{A}_2 be S-rings over the groups U and H/L respectively such that S is both an \mathcal{A}_1 - and an \mathcal{A}_2 -section with

$$(\mathcal{A}_1)_S = (\mathcal{A}_2)_S.$$

Then the set of all \mathcal{A} such that $\mathcal{A}_U = \mathcal{A}_1$ and $\mathcal{A}_{H/L} = \mathcal{A}_2$ has the smallest element, and it is a unique S -wreath product in this set.

Tensor (direct) product of S-rings

Let \mathcal{S} and \mathcal{T} be partitions of the groups H and K , and $\mathcal{S} \times \mathcal{T}$ the partition of $H \times K$ the classes of which are $S \times T$ where $S \in \mathcal{S}, T \in \mathcal{T}$.

Proposition

Let $\mathcal{A} = \langle \underline{\mathcal{S}} \rangle$ and $\mathcal{B} = \langle \underline{\mathcal{T}} \rangle$ be S-rings over H and K . Then $\mathcal{S} \times \mathcal{T}$ is an S-partition of $H \times K$; the S-ring of it is $\mathcal{A} \otimes \mathcal{B}$.

Exercise:

prove that $(\mathcal{A} \otimes \mathcal{B})_{(1,K)} \cong_{\text{Cay}} \mathcal{B}$ and $(\mathcal{A} \otimes \mathcal{B})/(1, K) \cong_{\text{Cay}} \mathcal{A}$.

Leung-Man and Schur-Wielandt theorem

Theorem (Leung-Man, 1994)

Let \mathcal{A} be a circulant S-ring. Suppose that \mathcal{A} is not a generalized wreath product. Then

$$\mathcal{A} = \mathcal{A}_0 \otimes \mathcal{A}_1 \otimes \cdots \otimes \mathcal{A}_k,$$

where \mathcal{A}_0 is cyclotomic and $\text{rk}(\mathcal{A}_i) = 2$ for all $i = 1, \dots, k$.

Remark: the factors \mathcal{A}_i are circulant S-rings of pairwise coprime degrees.

Theorem (Schur-Wielandt, 1964)

Let H be an abelian group of composite order with cyclic Sylow p -subgroup for some prime divisor p of $|H|$. Then any S-ring over H is imprimitive.

Schur theorem

Theorem

Let \mathcal{A} be an arbitrary S-ring over an abelian group H . Then $x^{(m)} \in \mathcal{A}$ for each $m \in \mathbb{Z}$ coprime to $|H|$ and all $x \in \mathcal{A}$.

Proof

- We may assume that $x = \underline{S}$ where $S \in \mathcal{S}$, and m is a prime;
- then $\underline{S}^m \equiv \underline{S}^{(m)} \pmod{m}$;
- since $\gcd(m, |H|) = 1$, this implies that $|S^{(m)}| = |S|$.
- Thus $S^{(m)} \in \mathcal{S}$.

Consequences of Schur's Theorem

Corollary 1

X is a basic set of \mathcal{A} if and only if $X^{(m)}$ is a basic set of \mathcal{A} .

Corollary 2

If \mathcal{A} is circulant S-ring over a group H , then $\text{Iso}_{\text{cay}}(\mathcal{A}) = \text{Aut}(H)$.

In particular, if \mathcal{A} is circulant, then the group $\text{Rad}(\mathcal{A}) := \text{Rad}(X)$ does not depend on $X \in \mathcal{S}$ containing a generator of H ; it is called the **radical** of \mathcal{A} .

Corollary 3

Any S-ring \mathcal{A} over a cyclic group C of prime order p is cyclotomic, that is $\mathcal{S} = \text{Orb}(K, C)$ for some $K \leq \text{Aut}(C)$.