## The technique of definable terms in Boolean valued analysis

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Let $\Phi$ be a set of first-order formulas of set-theoretic signature. A formula $\varphi$ is said to be of class $\Phi$ (" $\varphi$ is $\Phi$ " for short) whenever ZFC $\vdash\left[\varphi \Leftrightarrow \varphi^{\prime}\right]$ for some $\varphi^{\prime}$ in $\Phi$. Let $\tau(\bar{x})$ be any term introduced in (a conservative extension of) ZFC by means of a definition of the form $\tau(\bar{x})=y \Leftrightarrow \varphi(\bar{x}, y)$. Say that $\tau$ is of class $\Phi$ (" $\tau$ is $\Phi$ ") whenever $\varphi$ is of class $\Phi$. Say that $\tau$ is $\Phi$-definable via a term $\sigma$ (" $\tau$ is $\Phi(\sigma)$ ") whenever there is a formula $\varphi(\bar{x}, y, z)$ of class $\Phi$ such that ZFC $\vdash[\tau(\bar{x})=y \Leftrightarrow \varphi(\bar{x}, y, \sigma(\bar{x}))]$.

In what follows, we denote formulas and terms by $\varphi$ and $\tau, \sigma, \rho$ with possible indices; $\Delta_{0}$ is the smallest set containing the formulas $x \in y$ and closed under the connectives $\vee$, $\neg,(\exists x \in y) ; \quad \Sigma_{1}$ is constituted by the formulas $(\exists x) \varphi$, with $\varphi$ in $\Delta_{0}$. A formula $\varphi$ is of class $\Delta_{1}$ (" $\varphi$ is $\Delta_{1}$ ") whenever $\varphi$ and $\neg \varphi$ are $\Sigma_{1}$.

Lemma. (1) If $\varphi, \tau, \tau_{1}, \ldots, \tau_{n}$ are $\Sigma_{1}$ then so are $\varphi\left(\tau_{1}, \ldots, \tau_{n}\right)$ and $\tau\left(\tau_{1}, \ldots, \tau_{n}\right)$.
(2) If $\tau_{1}, \ldots, \tau_{n}$ are $\Sigma_{1}$ and $\varphi$ is $\Delta_{1}$ then $\varphi\left(\tau_{1}, \ldots, \tau_{n}\right)$ is $\Delta_{1}$.
(3) If $\tau$ is $\Sigma_{1}$ and $\varphi$ is $\Delta_{1}$ then $\{\tau(\bar{x}): \bar{x} \in y, \varphi(\bar{x}, y)\}$ is $\Sigma_{1}$.
(4) If $\tau$ is $\Sigma_{1}(\sigma)$ and $\rho$ is $\Sigma_{1}(\tau)$ then $\rho$ is $\Sigma_{1}(\sigma)$.
(5) If $\tau, \tau_{1}, \ldots, \tau_{n}$ are $\Sigma_{1}(\sigma)$ then so is $\tau\left(\tau_{1}, \ldots, \tau_{n}\right)$.
(6) If $\tau$ is $\Sigma_{1}$ then $\tau(\sigma)$ is $\Sigma_{1}(\sigma)$.
(7) If $\tau$ is $\Sigma_{1}$ and $\varphi$ is $\Delta_{1}$ then $\{\tau(\bar{x}): \bar{x} \in \sigma, \varphi(\bar{x}, \sigma)\}$ is $\Sigma_{1}(\sigma)$.
(8) If $\tau$ is $\Sigma_{1}$ and $\varphi$ is $\Delta_{1}$ then $\{\tau(\bar{x}): \bar{x} \in \sigma, \varphi(\bar{x}, \sigma)\}^{\mathbb{N}}$ is $\Sigma_{1}\left(\sigma^{\mathbb{N}}\right)$.

The following example shows that statements (3) and (7) do not extend to the case in which $\varphi$ is $\Sigma_{1}$.

Example. Assume that ZFC is consistent and put $\varphi(x):=(\exists z)(z \subseteq \mathbb{N} \wedge z \notin x)$. Then $\varphi$ is $\Sigma_{1}, \varphi$ is not $\Delta_{1}$, and $\{x \in y: \varphi(x)\}$ is not $\Sigma_{1}$.

In what follows, $(\cdot)^{\wedge}$ stands for the canonical embedding of $\mathbb{V}$ into the Boolean valued universe $\mathbb{V}^{(B)}$.

Theorem. If $\rho$ is $\Sigma_{1}, \tau$ is $\Sigma_{1}(\sigma)$, and all the parameters of $\rho, \sigma, \tau$ are in $\bar{x}$ then the following is provable in ZFC: for every complete Boolean algebra $B$ and all $\bar{x}$
(1) $\mathbb{V}^{(B)} \vDash\left[\rho(\bar{x})^{\wedge}=\rho\left(\bar{x}^{\wedge}\right)\right]$;
(2) $\mathbb{V}^{(B)} \vDash\left[\sigma(\bar{x})^{\wedge}=\sigma\left(\bar{x}^{\wedge}\right)\right] \Rightarrow \mathbb{V}^{(B)} \vDash\left[\tau(\bar{x})^{\wedge}=\tau\left(\bar{x}^{\wedge}\right)\right]$.

Let $\mathbb{R}_{\mathrm{D}}$ and $\mathbb{R}_{\mathrm{C}}$ stand for the set of reals defined as Dedekind cuts and, respectively, classes of Cauchy sequences in $\mathbb{Q}$.

Corollary (ZFC). Let $B$ be a complete Boolean algebra.
(1) $\mathbb{V}^{(B)} \vDash\left[\mathbb{R}_{\mathrm{D}} \subseteq \mathbb{R}_{\mathrm{D}}\right] ; \mathbb{V}^{(B)} \vDash\left[\mathcal{P}_{\mathrm{fin}}(X)^{\wedge}=\mathcal{P}_{\text {fin }}\left(X^{\wedge}\right)\right]$ for all $X$.
(2) The following properties of $B$ are pairwise equivalent: $B$ is $\sigma$-distributive; $\mathbb{V}^{(B)} \vDash\left[\mathcal{P}(\mathbb{N})^{\wedge}=\mathcal{P}(\mathbb{N})\right] ; \mathbb{V}^{(B)} \vDash\left[\left(\mathbb{N}^{\mathbb{N}}\right)^{\wedge}=\mathbb{N}^{\mathbb{N}}\right] ; \mathbb{V}^{(B)} \vDash\left[\mathbb{R}_{\mathrm{D}}^{\wedge}=\mathbb{R}_{\mathrm{D}}\right] ; \mathbb{V}^{(B)} \vDash\left[\mathbb{R}_{\mathrm{C}}^{\wedge} \subseteq \mathbb{R}_{\mathrm{C}}\right] ;$ $\mathbb{V}^{(B)} \vDash\left[\mathbb{R}^{\wedge}\right.$ and $\mathbb{R}$ are isomorphic ordered fields $]$.
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